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(12) **United States Patent**
Corts et al.

(10) **Patent No.:** **US 9,094,186 B2**
(45) **Date of Patent:** **Jul. 28, 2015**

(54) **SYSTEM AND METHOD FOR
TRANSMITTING DIGITAL MULTIMEDIA
DATA WITH ANALOG BROADCAST DATA**

(58) **Field of Classification Search**
CPC H04B 1/707; H04J 13/10; H04J 13/0022
USPC 375/219, 295; 705/40, 14
See application file for complete search history.

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(73) Assignee: **Impulse Radio, Inc.**, Larchmont, NY (US)

(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35 U.S.C. 154(b) by 0 days.

(21) Appl. No.: **13/795,229**

(22) Filed: **Mar. 12, 2013**

(65) **Prior Publication Data**

US 2015/0146711 A1 May 28, 2015

Related U.S. Application Data

(63) Continuation of application No. 10/484,518, filed as application No. PCT/US02/22898 on Jul. 17, 2002, now Pat. No. 8,396,100, said application No. 10/484,518 is a continuation-in-part of application No.

(Continued)

(51) **Int. Cl.**
H04B 17/02 (2006.01)
H04L 5/16 (2006.01)
H04L 7/04 (2006.01)

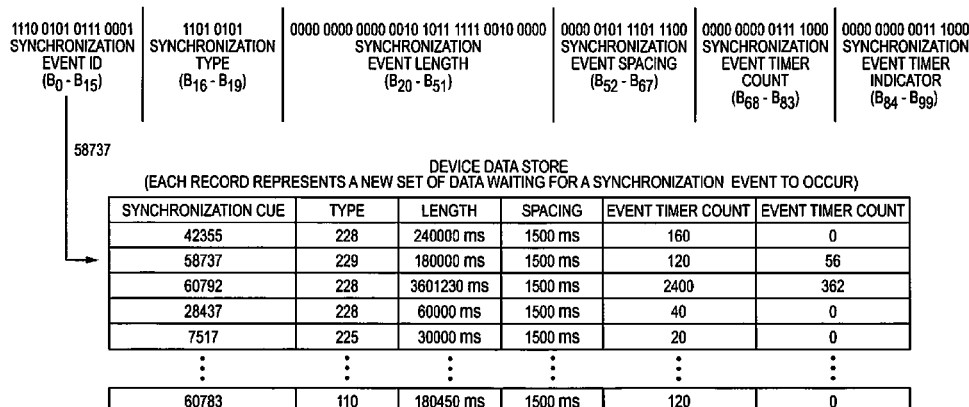
(52) **U.S. Cl.**
CPC **H04L 7/04** (2013.01)

(57) **ABSTRACT**

A method and system for the transmission of digital data (210) over existing analog radio frequencies (230) is presented, wherein the digital data may include audio data, visual data or audio-visual data for presentation either with analog broadcast data or at a selectable time. The digital data may be transmitted over a plurality of sub-channels that have varying degrees or reliability (250). A "quality-of-service" process manages the transmission of digital data over various sub-channels based on the reliability of the sub-channel, the amount of digital data and the type of digital data to be transmitted. The digital data may further be encrypted and authenticated.

18 Claims, 39 Drawing Sheets

DATA SYNCHRONIZATION STRUCTURE



Related U.S. Application Data

09/839,451, filed on Apr. 20, 2001, now Pat. No. 7,908,172, which is a continuation-in-part of application No. 09/802,469, filed on Mar. 9, 2001, now abandoned.

- (60) Provisional application No. 60/306,080, filed on Jul. 17, 2001, provisional application No. 60/188,050, filed on Mar. 9, 2000.

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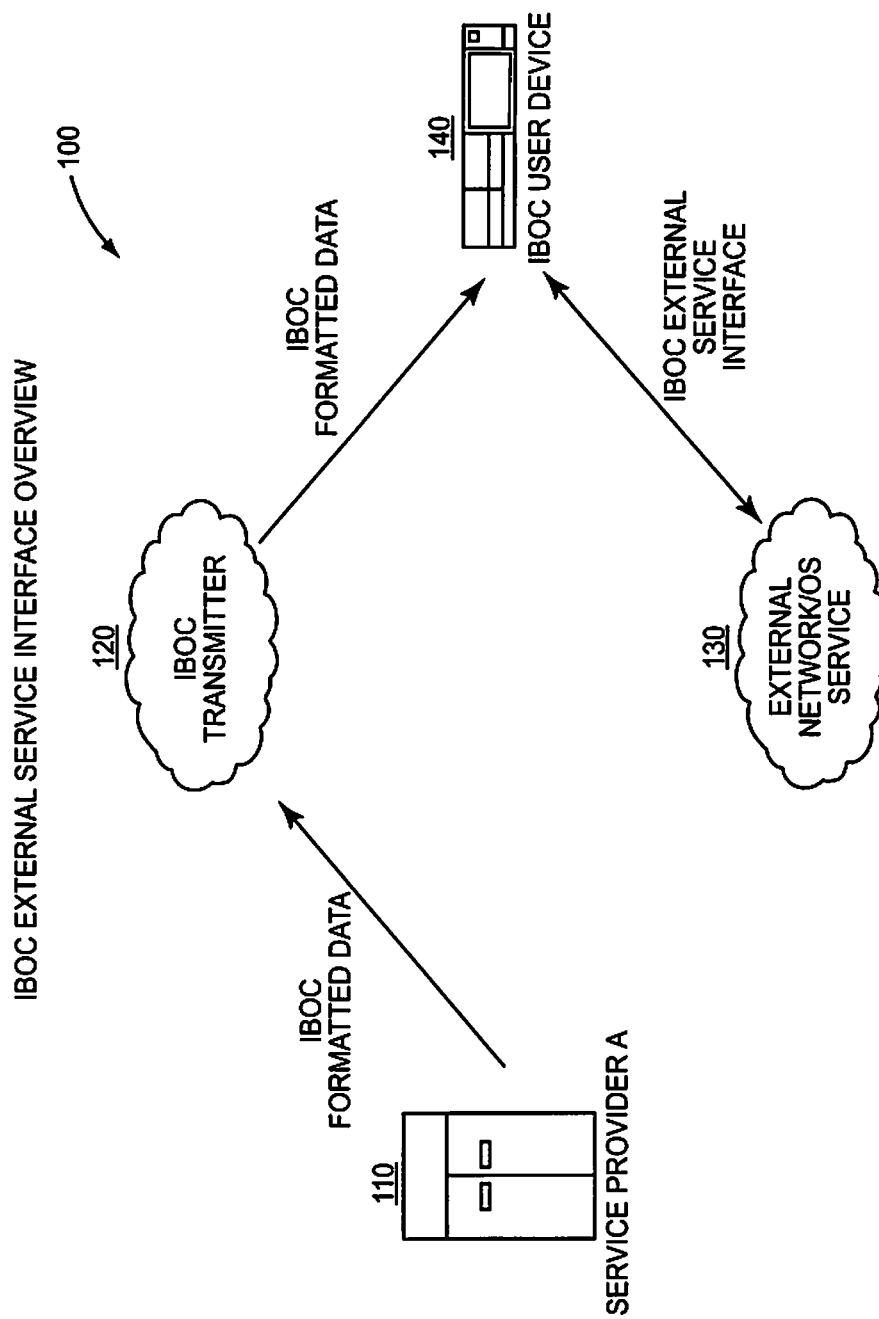


FIG. 1

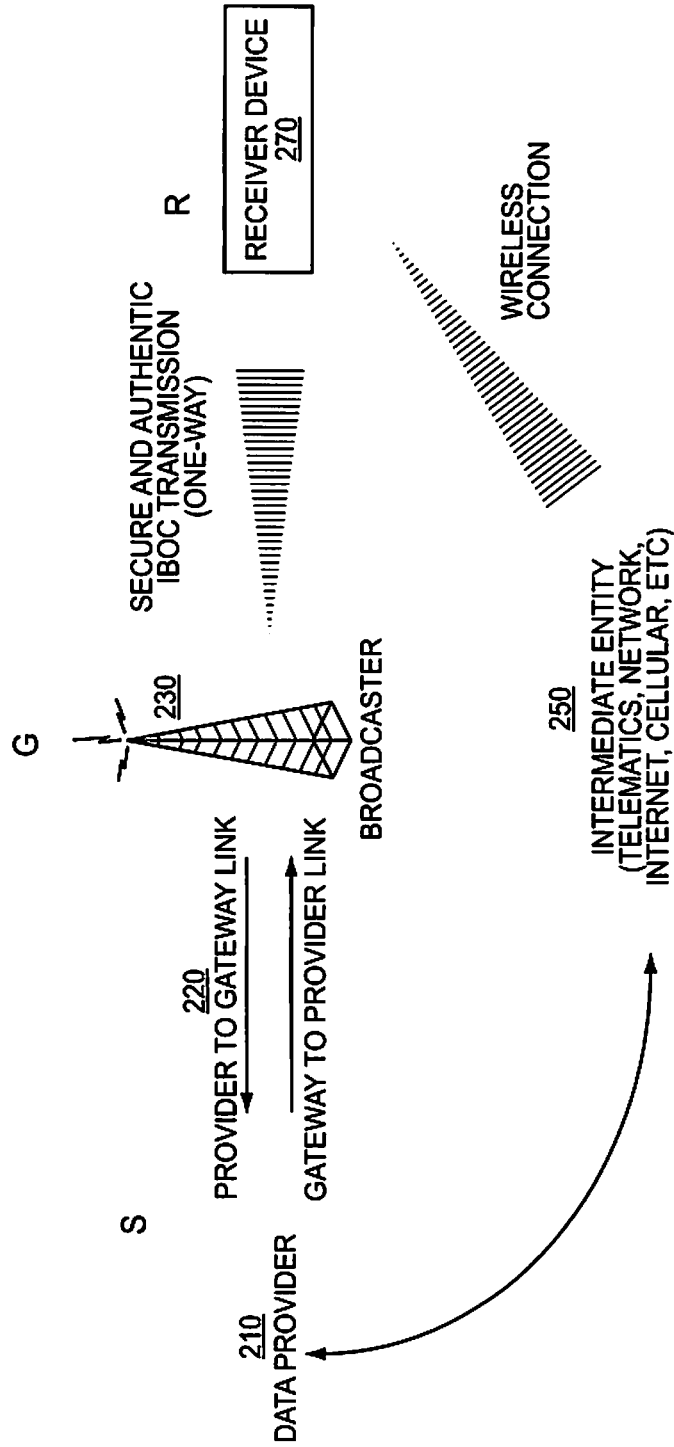
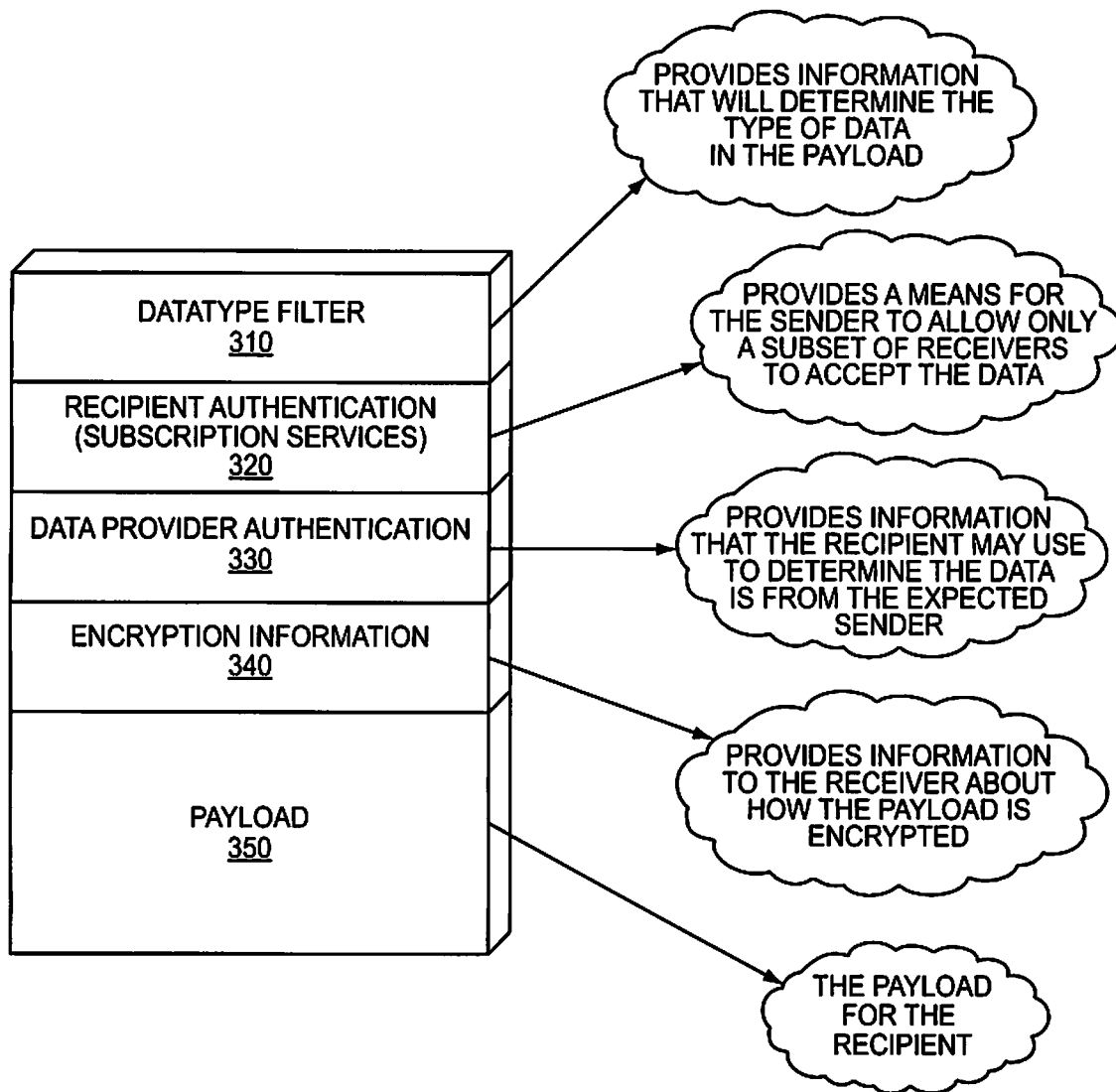


FIG. 2

**FIG. 3**

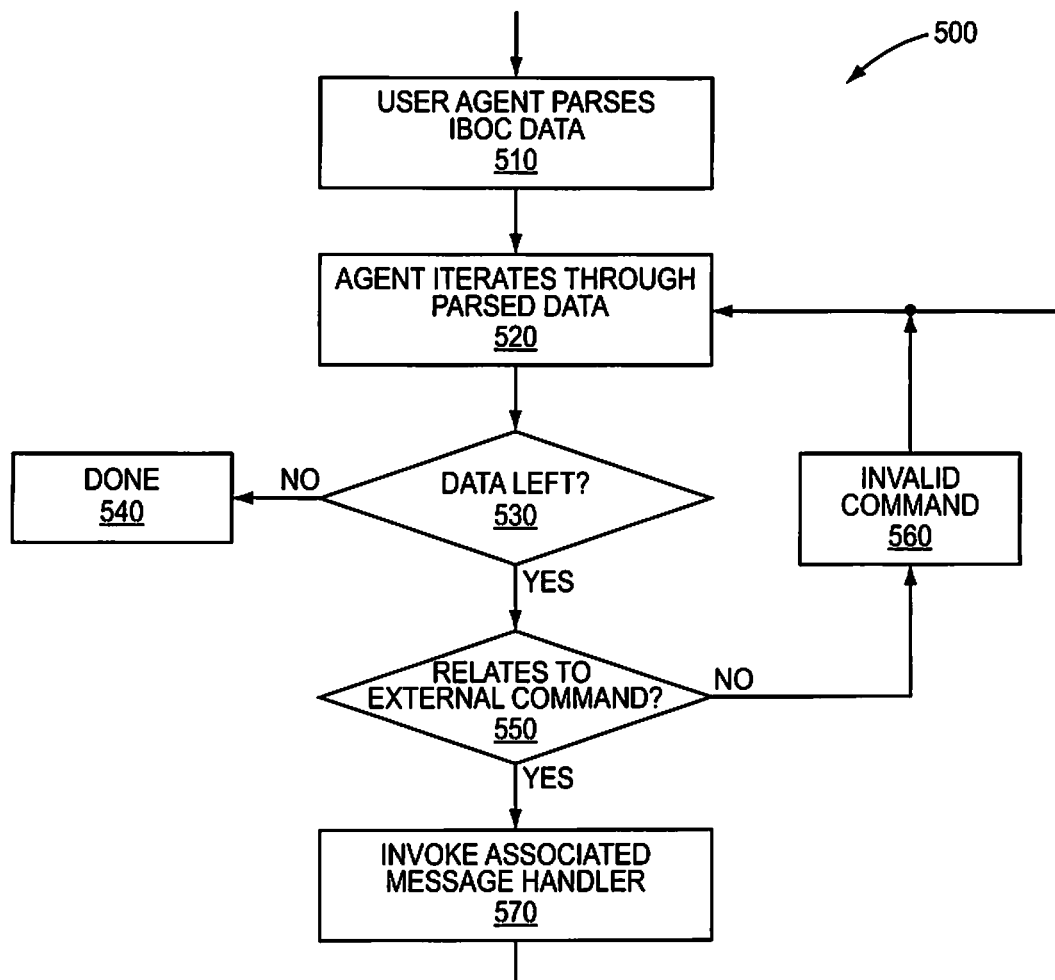
DATA SERVICE INFORMATION

DESCRIPTION	BIT POSITION	ABSTRACT
SYNCHRONIZATION	0 - 15	A 16 BIT WORD THAT MAY BE USED TO CALL THE CUE THE DATA ON DEMAND
SENDER TIME STAMP	16 - 47	A 32 BIT WORD REPRESENTING THE TIME THE DATA WAS SENT FROM THE BROADCASTER. THIS TIME MAY BE REPRESENTED IN MILLISECONDS.
RECEIVER TIME STAMP	48 - 79	A 32 BIT WORD REPRESENTING THE TIME LAST TIME DATA WAS RENDERED/EXECUTED BY THE RECEIVER.
DOMAIN ID	80-111	4 BYTES THAT MAY BE USED BY THE BROADCASTER TO IDENTIFY THE A DOMAIN ID OF THE SENDER. THE DATA SHOULD BE ABLE TO TRANSLATE TO THE EQUIVALENT OF 4 CALL LETTERS.
CONTENT RATING	112 - 115	BITS THAT MAY BE USED TO CONTENT RATING, E.G. 0001 - ALL AUDIENCES THRU 1111 - ADULTS ONLY
CONTENT CATEGORY	116 - 154	5 GROUPS OF 8 BITS THAT MAY BE USED TO CATEGORIZE THE DATA E.G., 00001000-SPORTS; 00000010 - MLB; 00000001 - YANKEES, 00000000 - N/A 00000000 - N/A
FILESIZE NUMBER	155 - 170	A 16 BIT WORD THAT MAY BE USED TO DETERMINE THE MAGNITUDE OF THE FILE SIZE BITS, (00 - BITS, 01 - BYTES, 10 - KILOBYTES, 11 - MEGABYTES)
FILESIZE MAGNITUDE	171 - 172	A 2 BIT WORD THAT MAY BE USED TO DETERMINE THE MAGNITUDE OF THE FILE SIZE BITS, (00 - BITS, 01 - BYTES, 10 - KILOBYTES, 11 - MEGABYTES)
RESERVED BIT	173 - 204	32 BIT RESERVED FOR FUTURE USE
STATUS BITS	205 - 236	32 INDIVIDUAL BITS THAT THE RECEIVER MAY USE TO AS STATUS FLAGS FOR THE DATA
EVENT ID	237 -252	16 BITS RESERVED TO IDENTIFY THE EVENT CUE MESSAGE USED TO SYNCHRONIZE THE DATA WITH THE AUDIO
EVENT INDICATOR	253 - 268	16 BITS USED TO IDENTIFY A POSITION WITHIN AN EVENT CUE MESSAGE THAT WILL START THIS USER DATA SHOULD BE PROCESSING
GROUP ID	269 - 284	16 BIT WORD THAT MAY BE USED TO GROUP THAT USER DATA BLOCK WITH THE DATA IN ANOTHER USER DATA BLOCK
MIME TYPE	285 - 300	16 BIT WORD THAT MAY BE USED TO IDENTIFY THE MIME TYPE OF THE USER DATA
USER DATA	301-	THE RENDERABLE/EXECUTABLE DATA

FIG. 4

SYNCHRONIZATION CUE 402
SENDER TIME STAMP 404
RECEIVER TIME STAMP 406
DOMAIN ID 408
CONTENT RATING 410
CONTENT CATEGORY LEVEL 1 THROUGH 5 412
FILE SIZE NUMBER 414
FILE SIZE MAGNITUDE 418
STATUS FLAGS RESERVED 420
STATUS FLAGS 422
EVENT ID 424
EVENT INDICATOR 426
GROUP ID 428
MIME TYPE 430
USER DATA 432

USER AGENT HANDLER INVOCATION

**FIG. 5**

USER AGENT HANDLER OPERATION

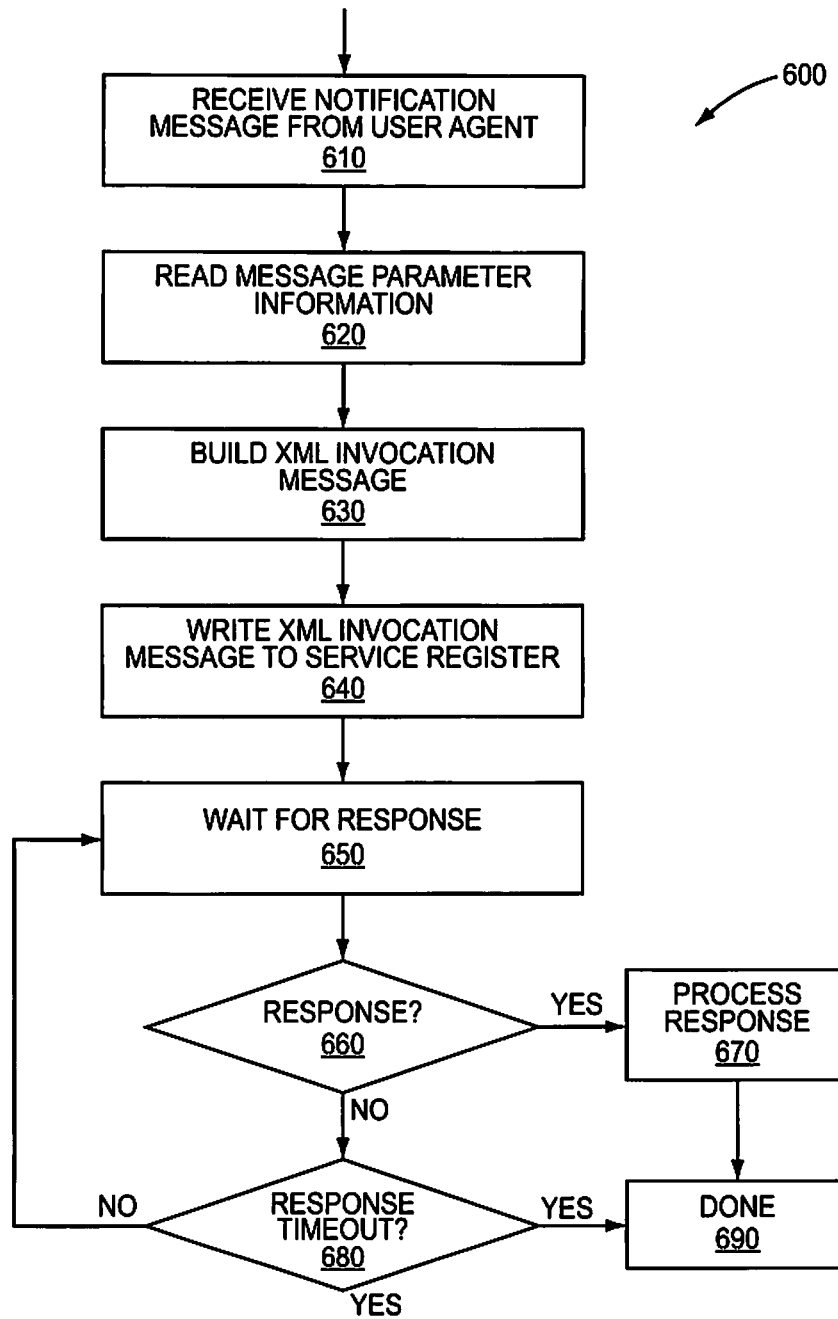
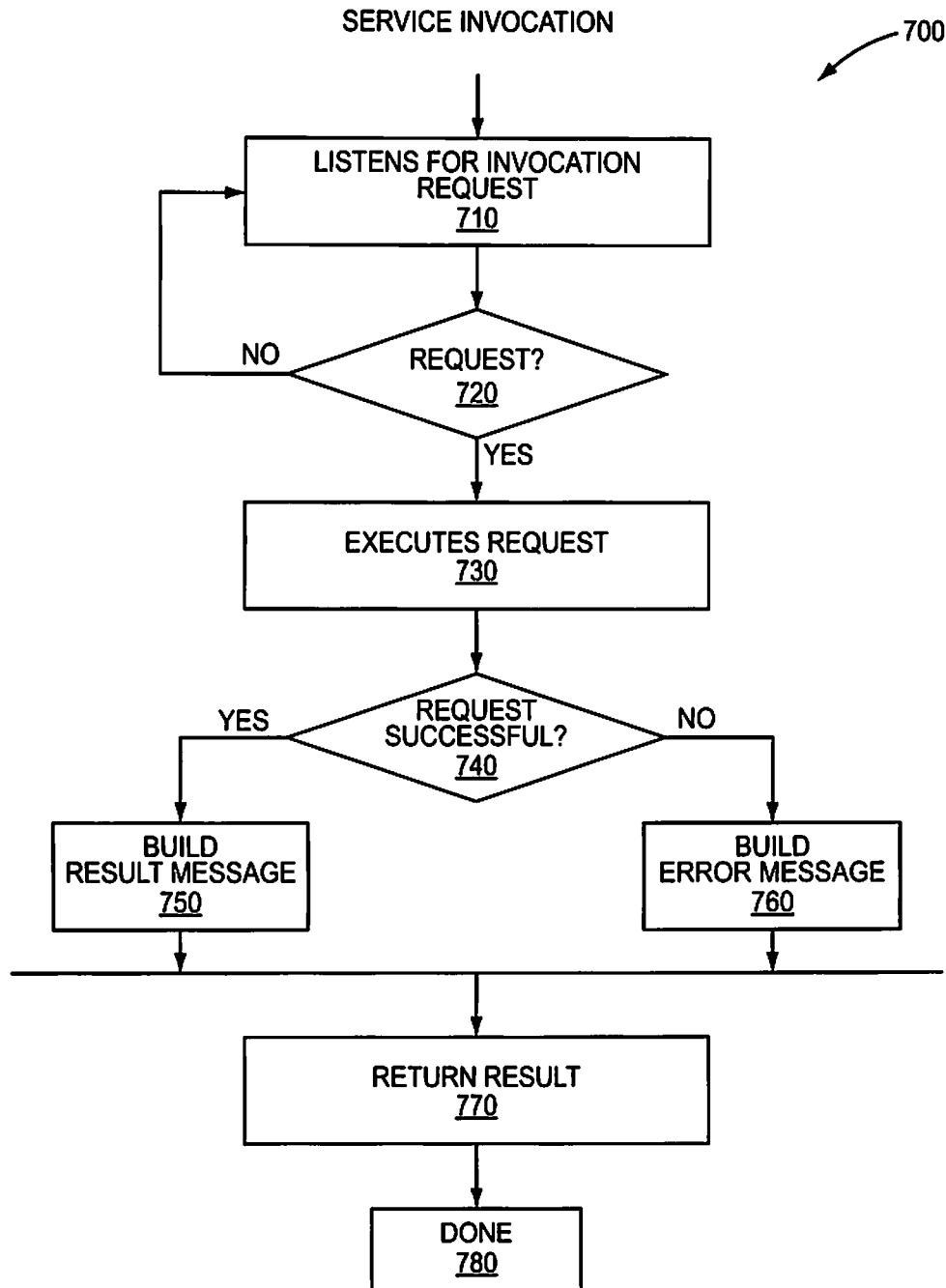
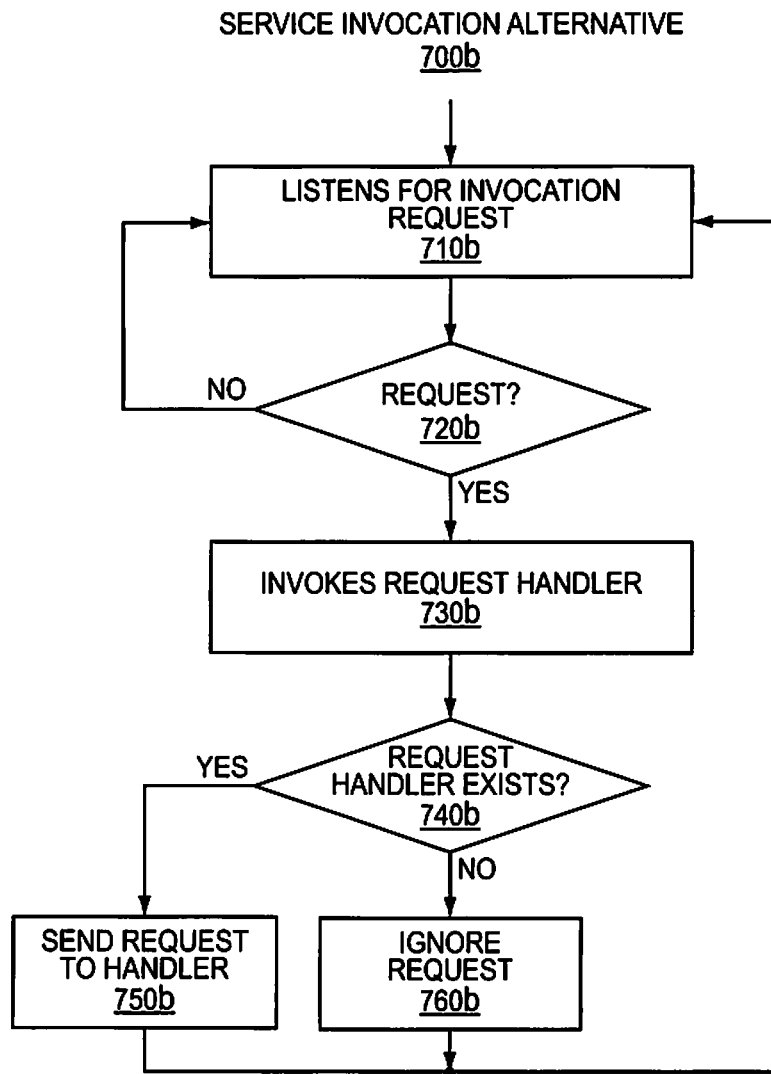
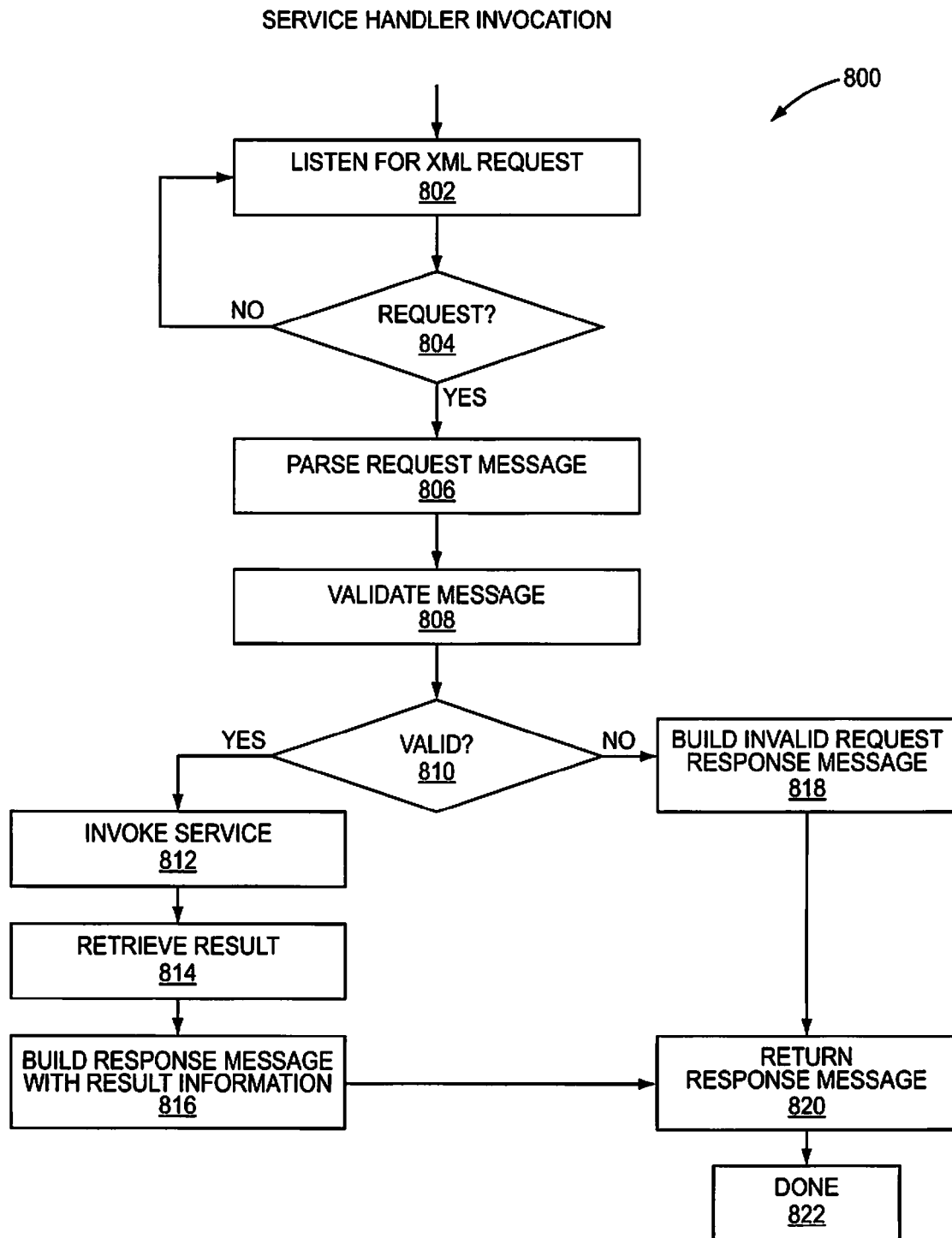


FIG. 6



**FIG. 7B**

**FIG. 8**

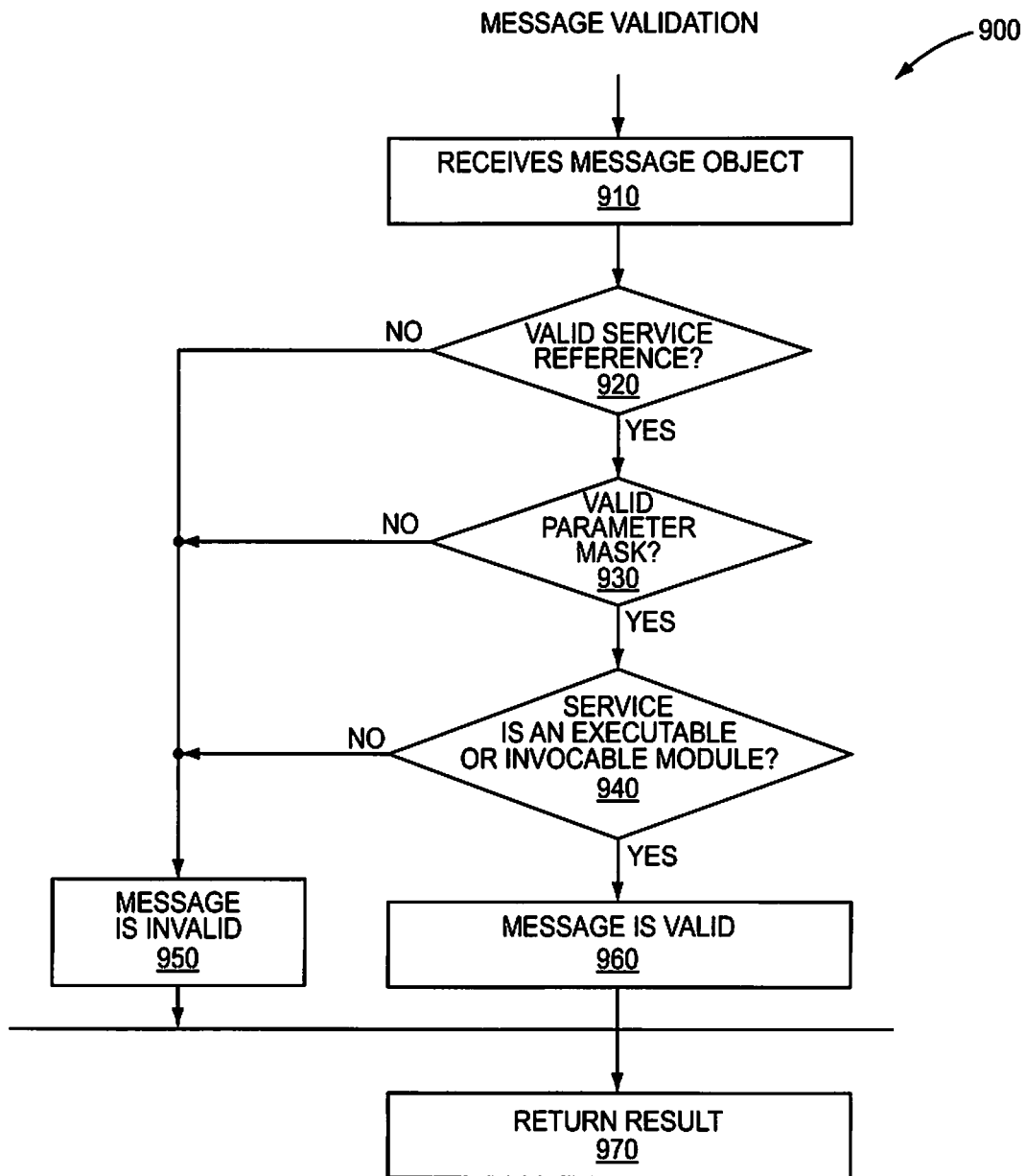


FIG. 9

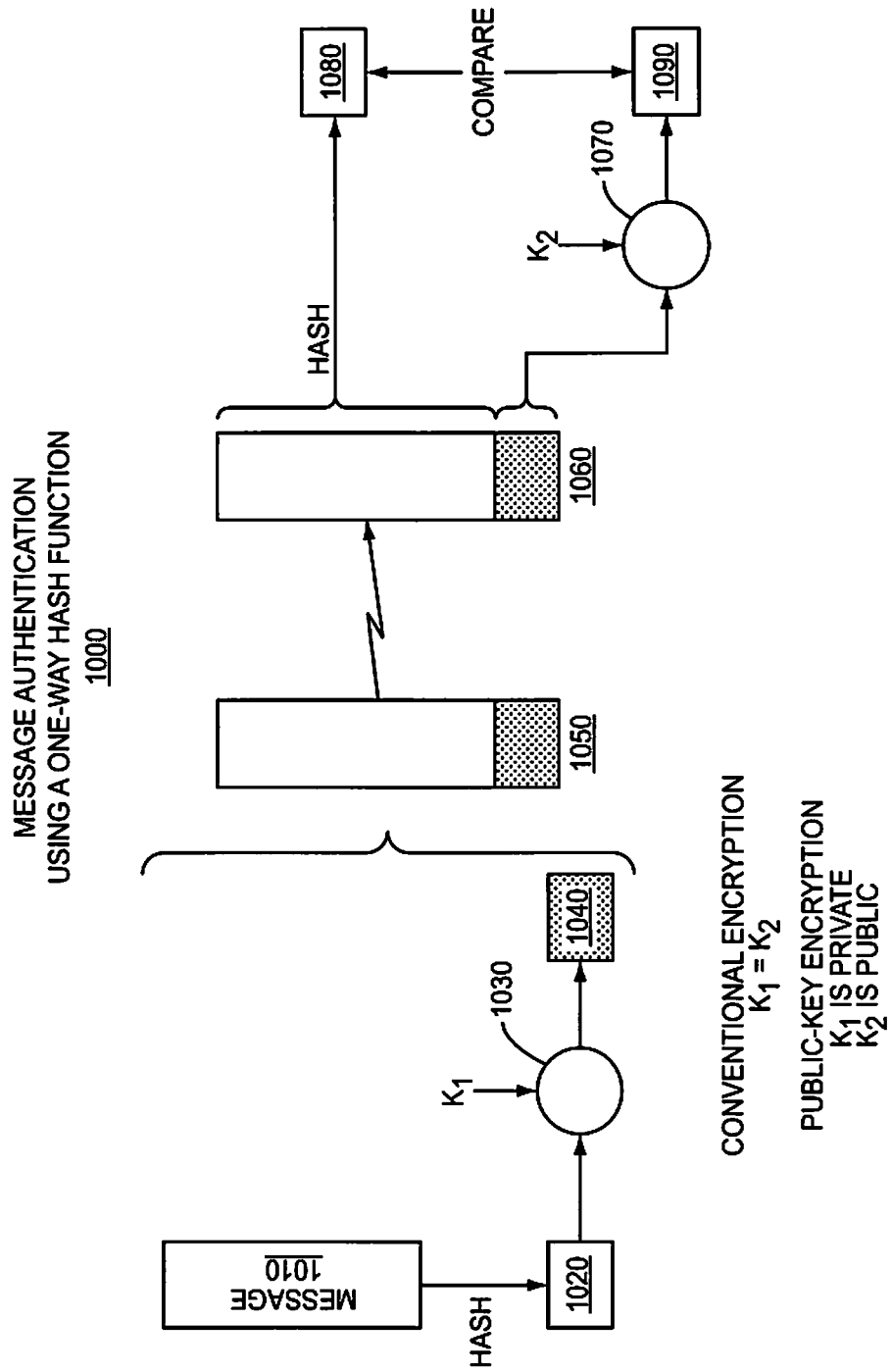


FIG. 10

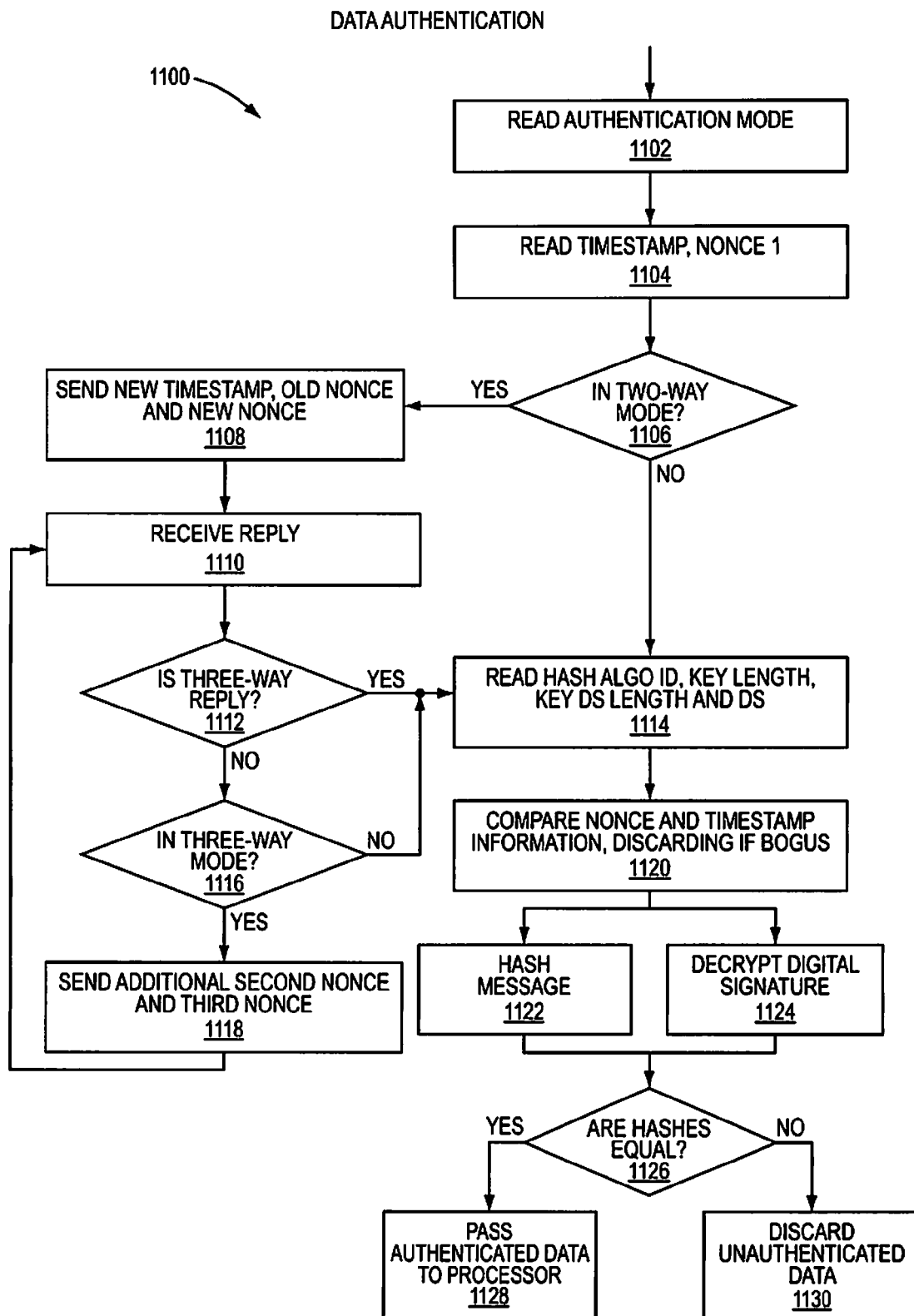
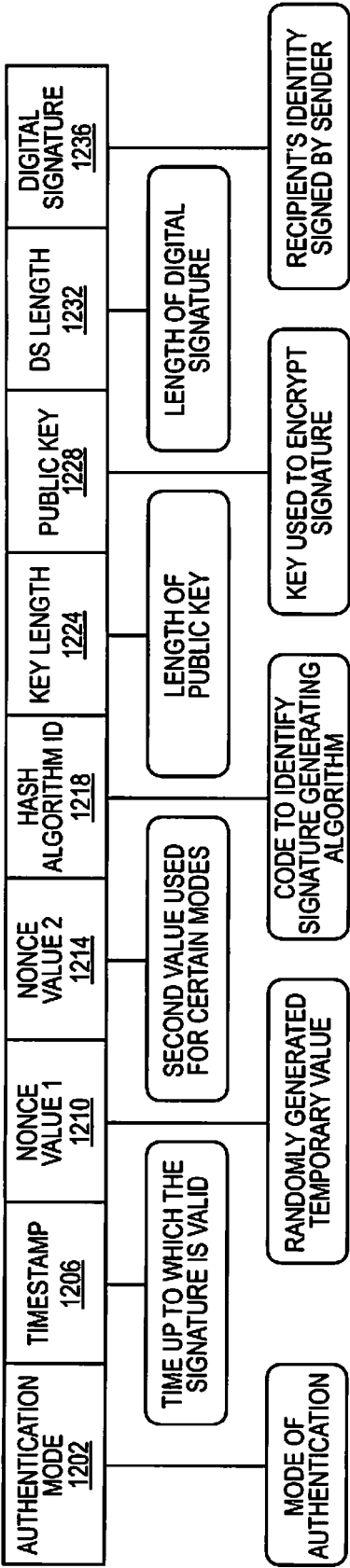
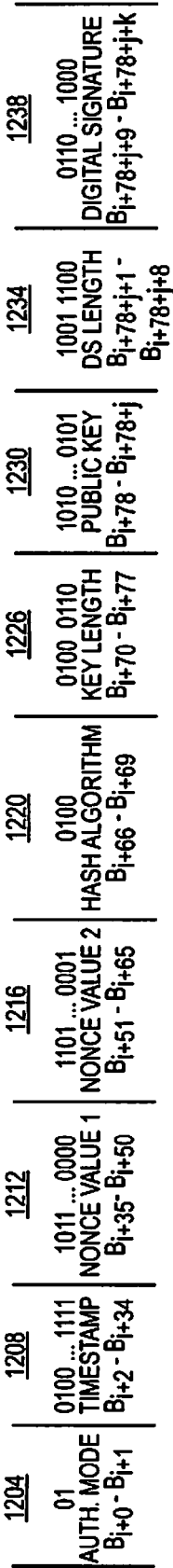


FIG. 11

AUTHENTICATION HEADER ROW



AUTHENTICATION HEADER EXAMPLE



j = PUBLIC KEY LENGTH IN BITS

k = DIGITAL SIGNATURE LENGTH IN BITS

FIG. 12

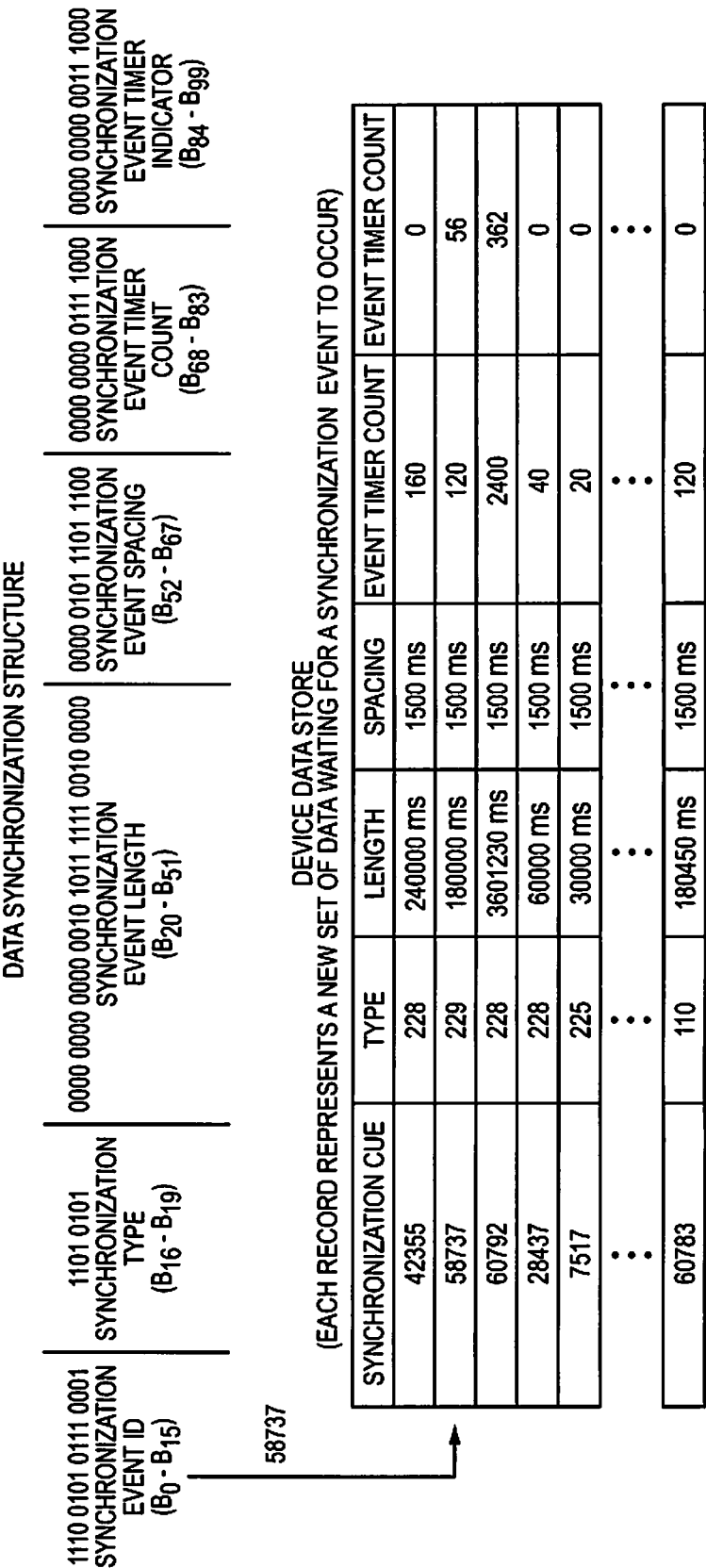
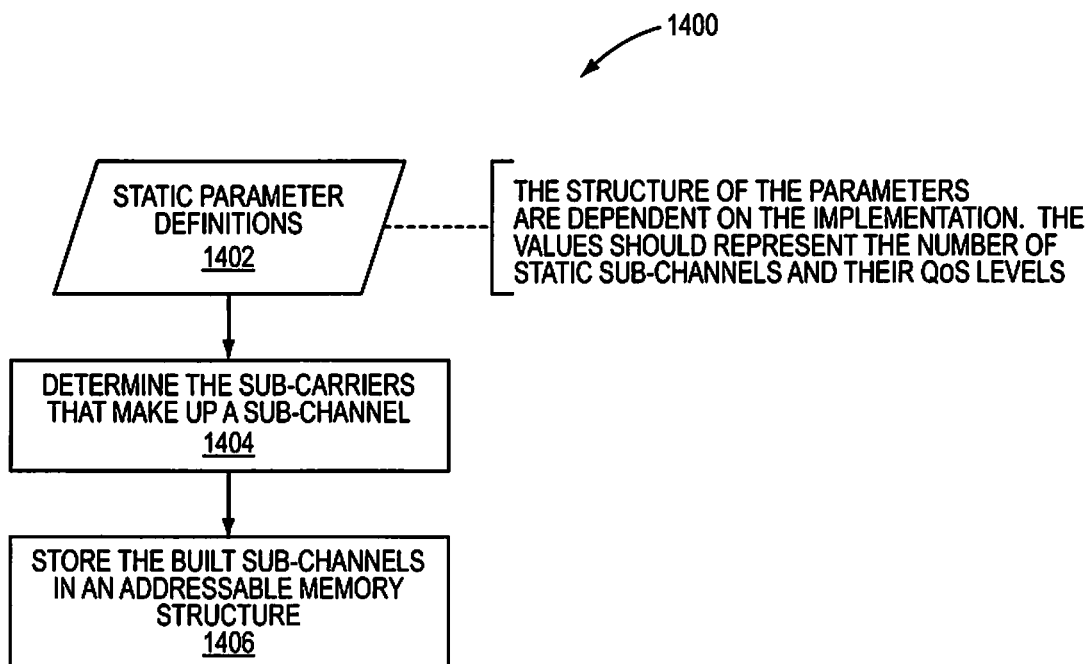


FIG. 13

**FIG. 14**

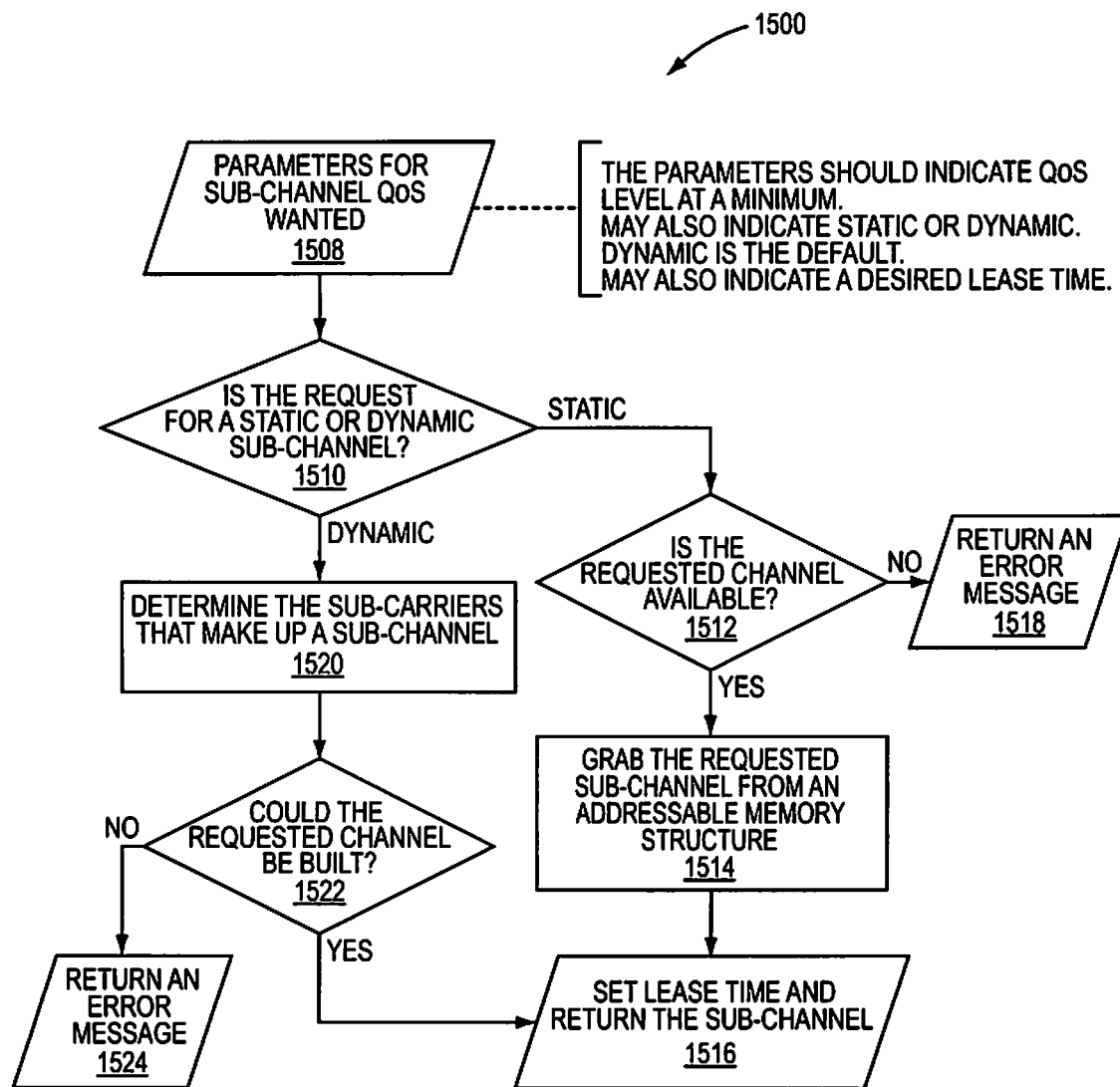


FIG. 15

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

Interface QoSManager 1610

public interface QoSManager

Handles the creation of subchannels and the level of service for a subchannel. It is also the point of entry for data services requiring a handle to sub-channels to send data

Method Summary

SubChannel	getSubChannel (int QoSLevel) 1620 Returns a subchannel of a specific QoSLevel returns null if the specific channel cannot be returned
------------	--

Method Detail

[getSubChannel](#) 1630

public [SubChannel](#) [getSubChannel](#) (int QoSLevel)
throws java.lang.Exception

Returns a subchannel of a specific QoSLevel returns null if the specific channel cannot be returned

Throws:

java.lang.Exception - if the Subchannel cannot be acquired.

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 16

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

Interface SubCarrier 1702

public interface SubCarrier

The SubCarrier class is the lowest level (aside from the physical layer) of an IBOC transmission. Multiple SubCarriers are grouped together to form a SubChannel. SubCarriers have a Quality of Service (QoS) rating that describes the reliability of the carrier. Generally, the higher the QoS rating, the lower the reliability of the carrier.

Method Summary	
void	clear () 1704 Returns the carrier to “pristine” condition for return to the pool.
int	getQoSRating () 1706 Gets the Quality of Service (QoS) rating of the SubCarrier
int	read (byte [] b, int len) 1708 Reads up to len bytes off of the carrier into byte-array b.
int	write (byte [] b) 1710 Writes up to b.length bytes on the carrier from byte-array b.

Method Detail

[getQoSRating](#) 1712

public int getQoSRating ()

Gets the Quality of Service (QoS) rating of the SubCarrier

Returns:

the carrier’s rating

[read](#) 1714

public int read(byte[] b,
int len)

FIG. 17A

Reads up to len bytes off of the carrier into byte-array b

Parameters:

b - array to be populated with the bytes read

len - number of bytes requested to be read

Returns:

the number of bytes actually read

write [1716](#)

public int write (byte[] b)

Writes up to b.length bytes on to the carrier from byte-array b.

Parameters:

b - an array of bytes to be written

Returns:

the number of bytes actually written

clear [1718](#)

public void clear ()

Returns the carrier to “pristine” condition for return to the pool.

Class [Tree](#) [Deprecated](#) [Index](#) [Help](#)

[PREV CLASS](#) [NEXT CLASS](#)

[FRAMES](#) [NO FRAMES](#)

[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)

[DETAIL FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 17B

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL FIELD](#) | [CONSTR](#) | [METHOD](#)

Interface SubCarrierFactory [1802](#)

public interface SubCarrierFactory

An instance of SubCarrierFactory is used to create SubCarriers that comprise a SubChannel.

Method Summary

SubCarrier	newSubCarrier (int rating) 1804 Create a new SubCarrier with the specified Quality of Service (QoS) rating.
------------	--

Method Detail

[newSubCarrier](#) [1806](#)

public [SubCarrier](#) newSubCarrier (int rating)

Create a new SubCarrier with the specified Quality of Service (QoS) rating.

Parameters:

rating - the level of reliability

Returns:

a SubCarrier object with the specified rating

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 18

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) [FIELD](#) [CONSTR](#) [METHOD](#)[DETAIL](#) [FIELD](#) [CONSTR](#) [METHOD](#)

Interface SubCarrierPool 1902

public interface SubCarrierPool

The interface by which subcarriers can be locked/unlocked for use by a subchannel. The storage of a predefined amount of subcarrier objects in this pool prevents unnecessary reconstruction and garbage collection. All implementing classes should provide at least one constructor with a single argument of type int and constructs this many SubCarrier objects to hold in a Collection.

Method Summary	
void	checkIn (SubCarrier subc) 1904 A method to notify the pool that subc is available for checkout.
SubCarrier	checkOut (int qosRating) 1906 Returns a Subcarrier object with a quality of service rating matching qosRating after removing it from the pool.
int[]	getAvailableRatings () 1908 Returns an array of ints matching the quality of service of each subcarrier currently in the pool
int	getCount () 1910 Returns the size of the pool

Method Detail

checkOut

public [SubCarrier](#) checkOut (int qosRating) 1912
throws java.lang.Exception

Returns a Subcarrier object with a quality of service rating matching qosRating after removing it from the pool

Parameters:

qosRating - the quality of service rating of the subcarrier

Throws:

java.lang.Exception - when a subcarrier with qosRating is unavailable

FIG. 19A

checkIn

public void checkIn (SubCarrier subc) 1914
throws java.lang.Exception

A method to notify the pool that subc is available for checkout

Parameters:

subc - the subcarrier ready to be checked in

Throws:

java.lang.Exception - when a subcarrier of the same quality of service exists in the pool already

getAvailable Ratings 1916

public int [] getAvailableRatings ()

Returns an array of ints matching the quality of service of each subcarrier currently in the pool.

getCount 1918

public int getCount ()

Returns the size of the pool.

Class [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 19B

DETAIL FIELD | CONSTR | METHOD

2002

Transport mechanism used by data services to send data. It is also use by receiving devices to read data sent by a data service.

Method Summary

```
public java.io.OutputStream getOutputStream ( )
    throws java.io.IOException
    Gets the output Stream used to send data to the channel
    Throws:
        java.io.IOException - if the OutputStream cannot be returned
```

```
public java.io.InputStream getInputStream ()
    throws java.io.IOException
    Gets the output Stream used to read data from the channel
```

FIG. 20A

Throws:

java.io.IOException - if the OutputStream cannot be returned

getQosLevel 2016

public int getQosLevel ()

Gets the Quality of Servicelevel of the channel

destroy 2018

public int () destroy ()

perform any cleanup work required to destroy the channel

Returns:

an array integers that determine the sub-carriers that compose this channel.

Class [Tree](#) [Deprecated](#) [Index](#) [Help](#)

[PREV CLASS](#) [NEXT CLASS](#)

[FRAMES](#) [NO FRAMES](#)

[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)

[DETAIL FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 20B

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

Interface SubChannelFactory 2102

public interface SubChannelFactory

Factory class for creating new subChannels

Method Summary

SubChannelFactory	newSubChannel (SubCarrier [] subcars) 2104 method for creating a new subchannel
-----------------------------------	--

Method Detail

newSubChannel

public [SubChannel](#) newSubChannel ([SubCarrier](#) [] subcars) 2106
throws java.lang.Exception

method for creating a new subchannel

Parameters:

subcars - an array of Subcarriers used to compose the new sub channel

Throws:

java.lang.Exception - if the Subchannel cannot be created

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 21

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) [FIELD](#) [CONSTR](#) [METHOD](#)[DETAIL](#) [FIELD](#) [CONSTR](#) [METHOD](#)**Interface Service** [2202](#)

public interface Service

Method Summary	
void	Authenticate (DeviceKey dKey) 2204
java.io.InputStream	getInputStream () 2206
java.io.OutputStream	getOutputStream () 2208
ServiceMetaData	getServiceMetaData () 2210
void	getServiceMetaData () 2212

Method Detail[getServiceMetaData](#) [2214](#)

public [ServiceMetaData](#) [getServiceMetaData \(\)](#)

[getServiceMetaData](#) [2216](#)

public void [getServiceMetaData \(\)](#)

[getOutputStream](#) [2218](#)

public java.io.OutputStream [getOutputStream \(\)](#)
throws java.io.IOException

Gets the output Stream used to read data from the channel

FIG. 22A

getInputStream 2220

```
public java.io.InputStream getInputStream ()
    throws java.io.IOException
```

Authenticate 2222

```
public void Authenticate (DeviceKey dKey)
                        throws java.lang.Exception
```

Class [Tree](#) [Deprecated](#) [Index](#) [Help](#)

PREV CLASS NEXT CLASS

FRAMES NO FRAMES

SUMMARY INNER | FIELD | CONSTR | METHOD

DETAIL FIELD | CONSTR | METHOD

FIG. 22B

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

Interface **ServiceListener** [2302](#)

public interface ServiceListener

Implementations of this interface are responsible for delegating the construction of subchannel objects and handing them off to the appropriate handler objects. This requires access to both the SubcarrierPool and the SubchannelFactory.

Method Summary

Service []	getServices () 2304 returns an array of the services that are actively being received
-----------------------------	---

Method Detail

[getServices](#) [2306](#)

public [Service](#)[] [getServices](#) ()
returns an array of the services that are actively being received

[Class](#) [Tree](#) [Deprecated](#) [Index](#) [Help](#)[PREV CLASS](#) [NEXT CLASS](#)[FRAMES](#) [NO FRAMES](#)[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 23

Class [Tree](#) [Deprecated](#) [Index](#) [Help](#)

[PREV CLASS](#) [NEXT CLASS](#)

[FRAMES](#) [NO FRAMES](#)

[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)

[DETAIL](#) [FIELD](#) | [CONSTR](#) | [METHOD](#)

Interface ServiceMetaData

[2402](#)

public interface ServiceMetaData

Method Summary	
int	getCategory (int level) 2404
int	getContentRating () 2406
int	getDataSize () 2408
int	getDataSizeMagnitude () 2410
byte[]	getDomainID () 2412
int	getEventIndicator () 2414
int	getGroupID () 2416
int	getMimeType () 2418
long	getReceiverTimeStamp () 2420
int	getReservedBits () 2422
long	getSenderTimeStamp () 2424
int	getStatusBits () 2426
int	getSyncCue () 2428
void	setReceiverTimeStamp (long tStamp) 2430

FIG. 24A

Method Detail

getSyncCue

```
public int getSyncCue ( )
```

getSenderTimeStamp

```
public long getSenderTimeStamp ( )
```

setReceiverTimeStamp

```
public void setReceiverTimeStamp (long tStamp)
```

getReceiverTimeStamp

```
public long getReceiverTimeStamp ( )
```

getDomainID

```
public byte [ ] getDomainID ( )
```

getContentRating

```
public int getContentRating ( )
```

getCategory

```
public int getCategory (int level)
```

getDataSize

```
public int getDataSize ( )
```

FIG. 24B

getDataSizeMagnitude

```
public int getDataSizeMagnitude ( )
```

getReservedBits

```
public int getReservedBits ( )
```

getStatusBits

```
public int getStatusBits ( )
```

getEventIndicator

```
public int getEventIndicator (int level)
```

getGroupID

```
public int getGroupID ( )
```

getMimeType

```
public int getMimeType ( )
```

Class [Tree](#) [Deprecated](#) [Index](#) [Help](#)

[PREV CLASS](#) [NEXT CLASS](#)

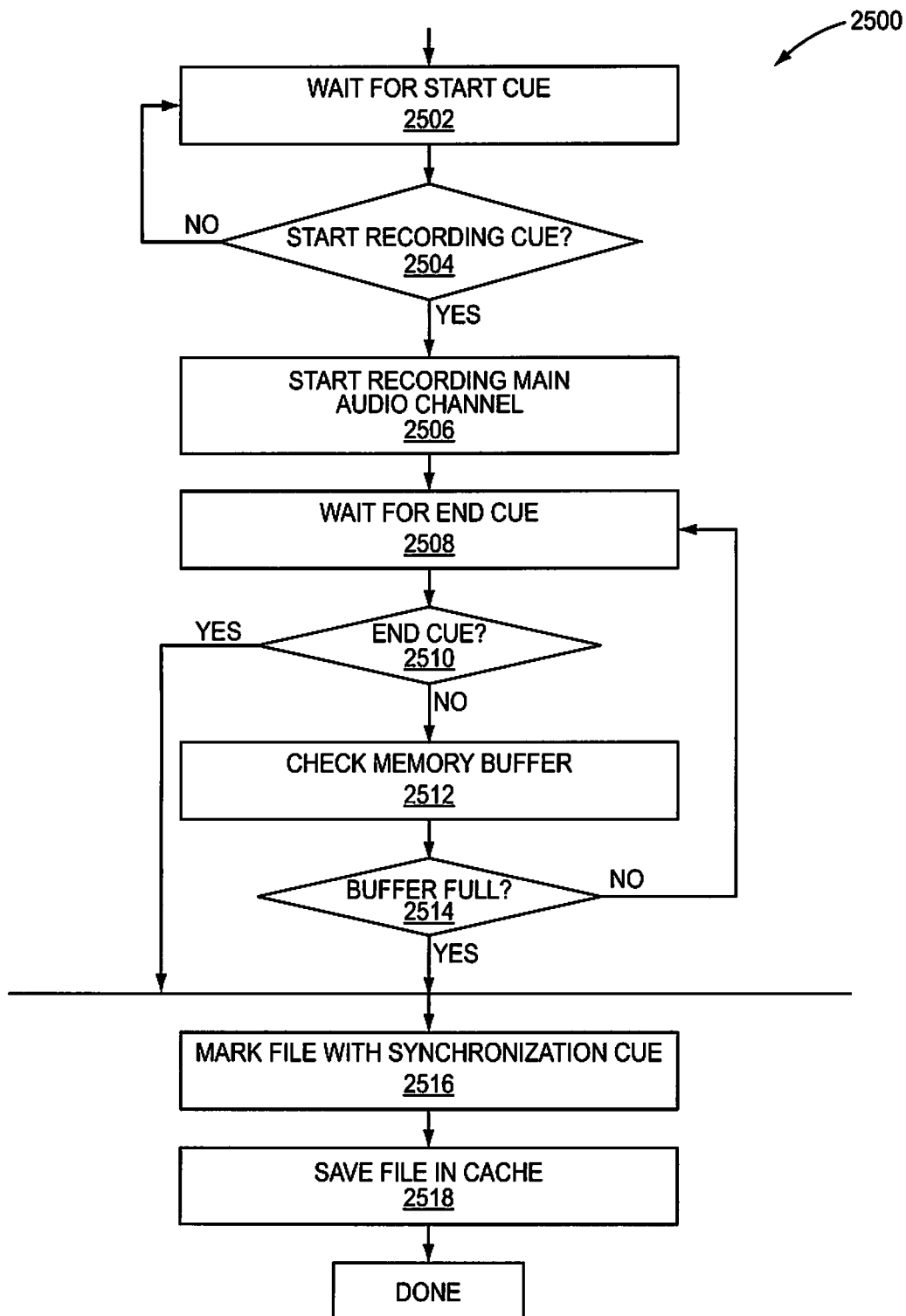
[FRAMES](#) [NO FRAMES](#)

[SUMMARY](#) [INNER](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)

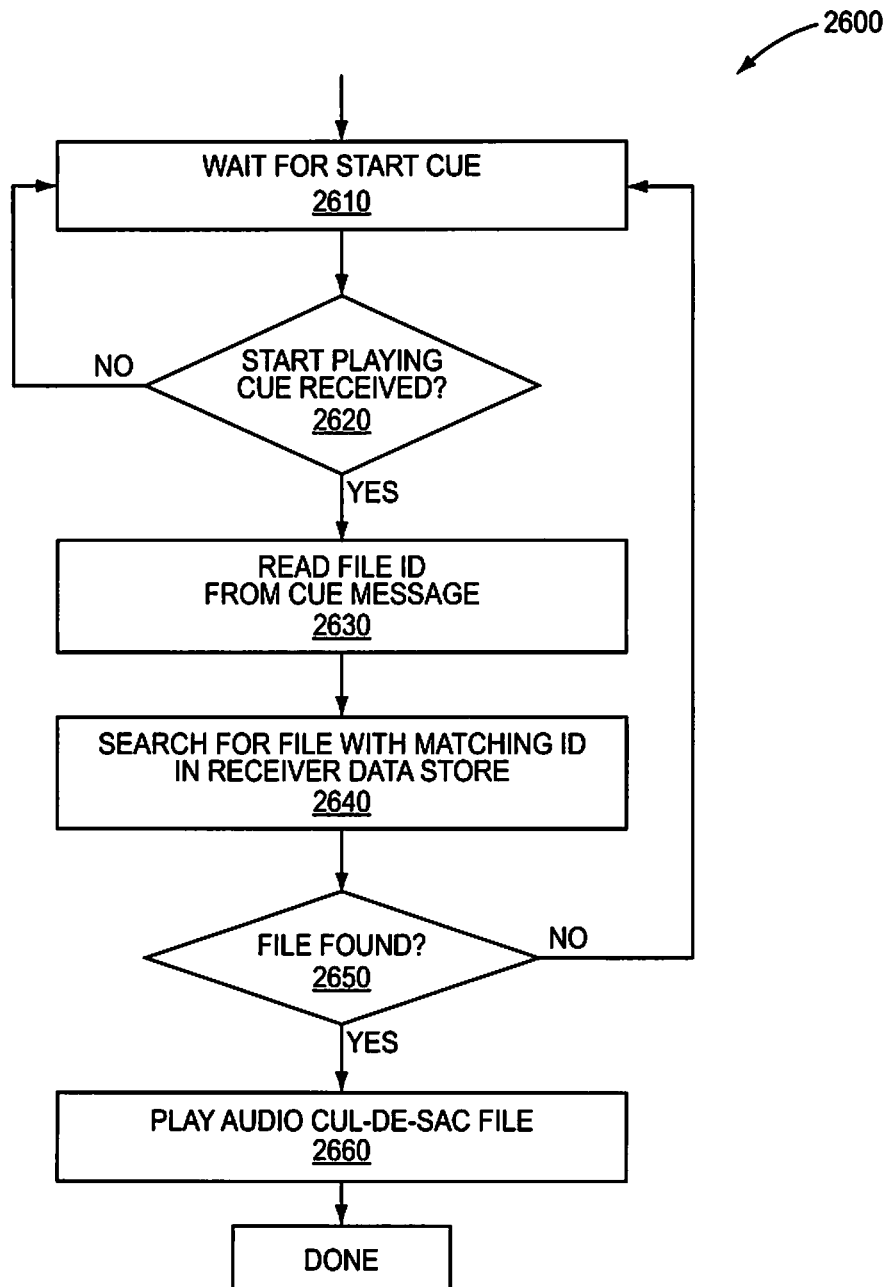
[DETAIL FIELD](#) | [CONSTR](#) | [METHOD](#)

FIG. 24C

USING SYNCHRONIZATION CUE TO CREATE AUDIO CUL-DE-SAC

**FIG. 25**

USING SYNCHRONIZATION CUE TO TRIGGER AN AUDIO CUL-DE-SAC

**FIG. 26**

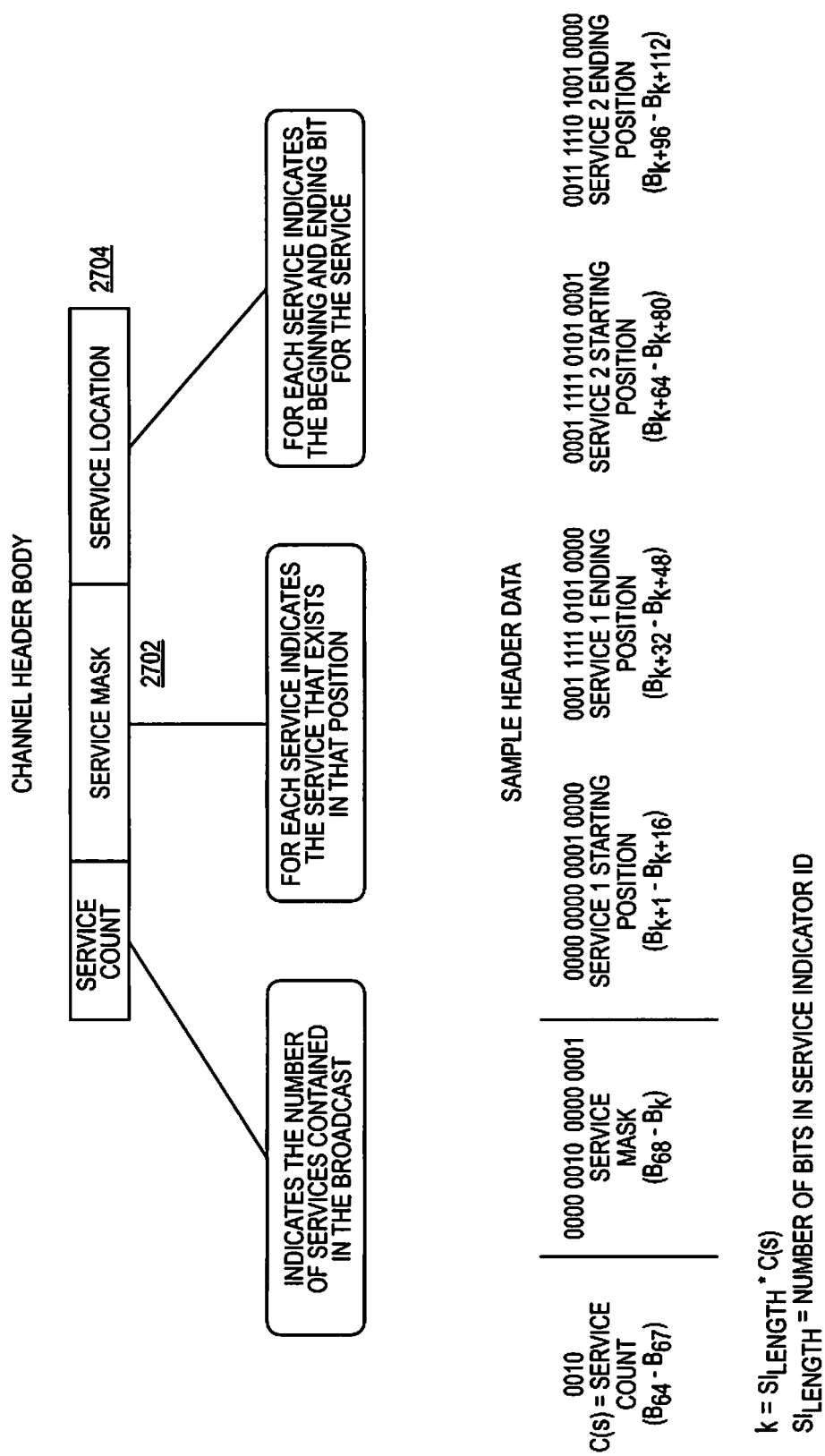


FIG. 27

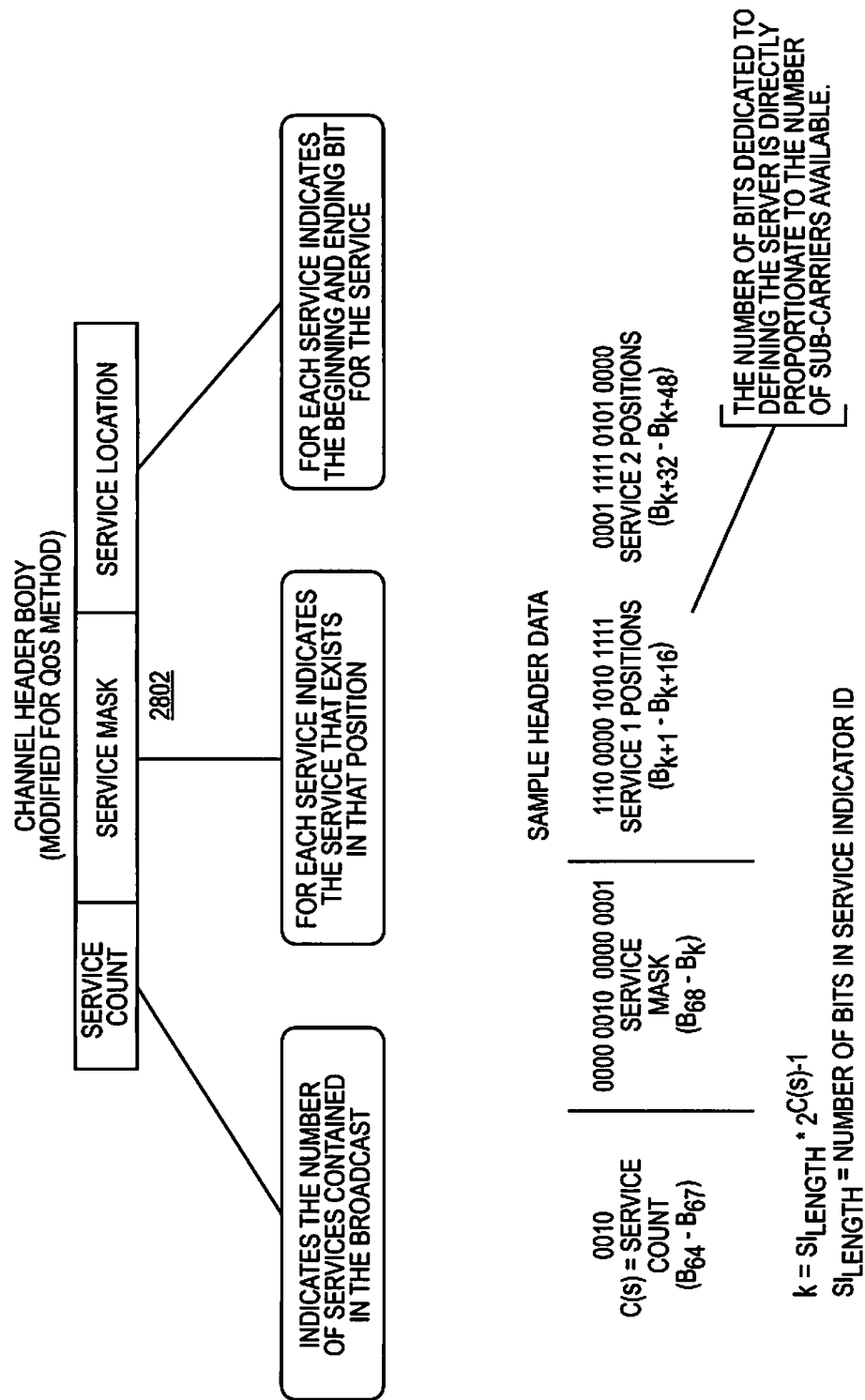
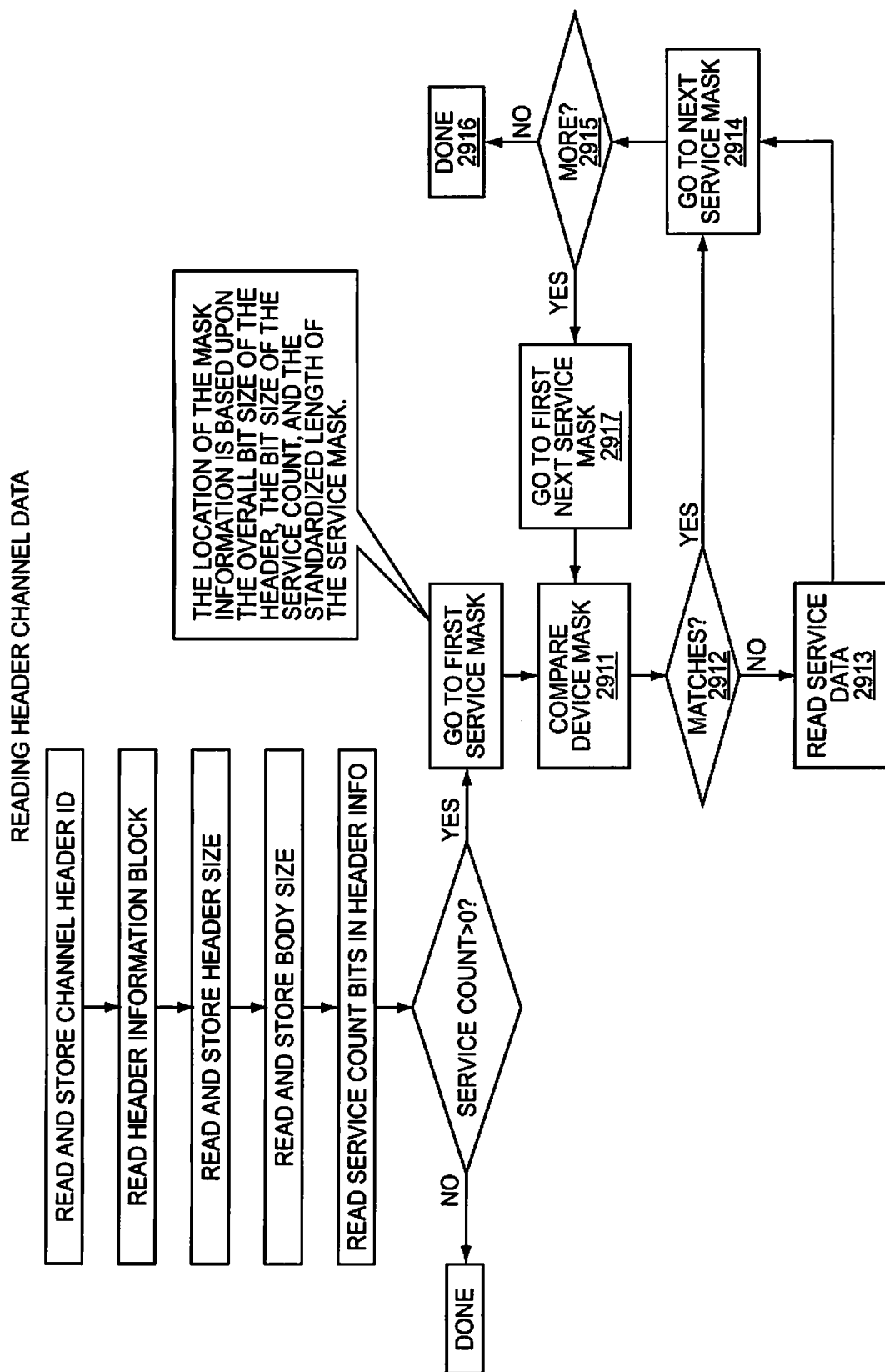


FIG. 28



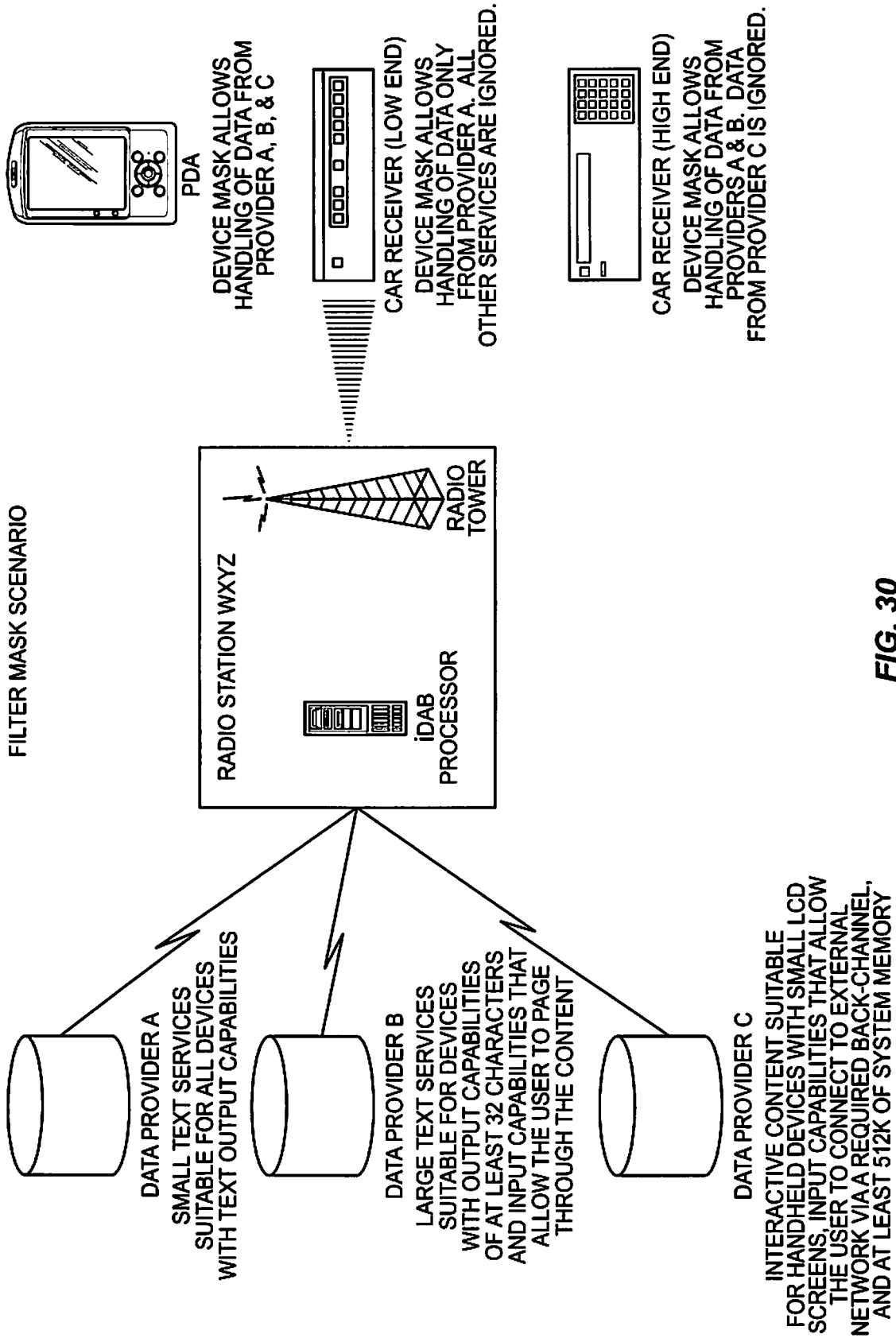


FIG. 30

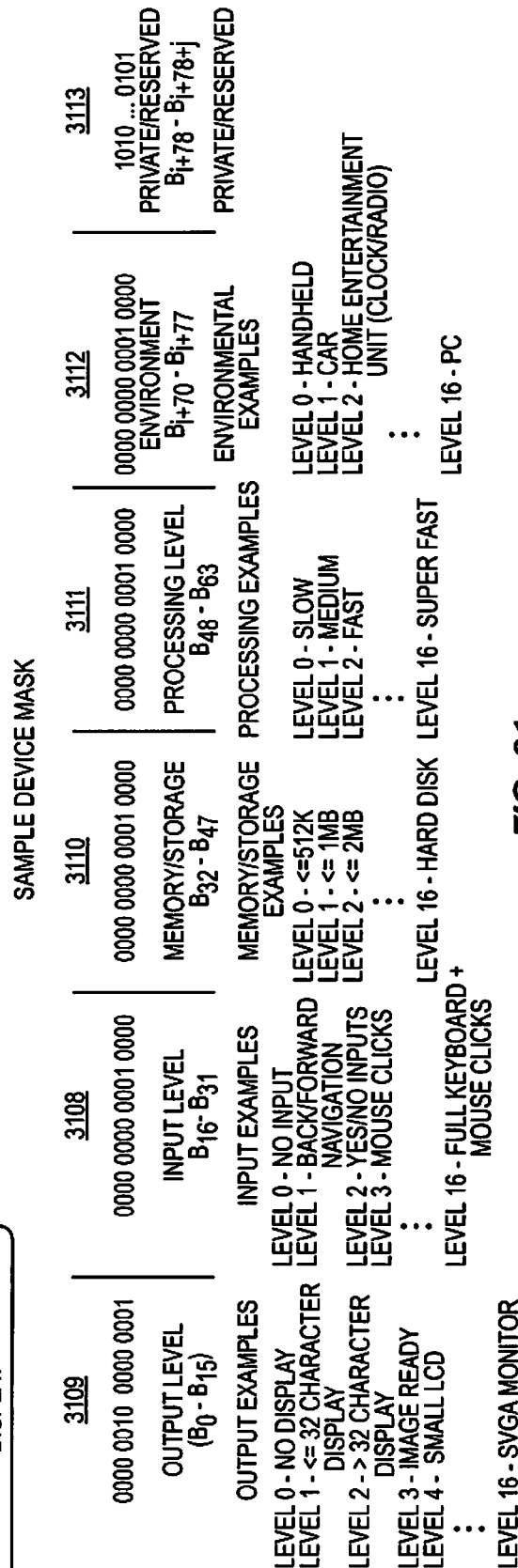
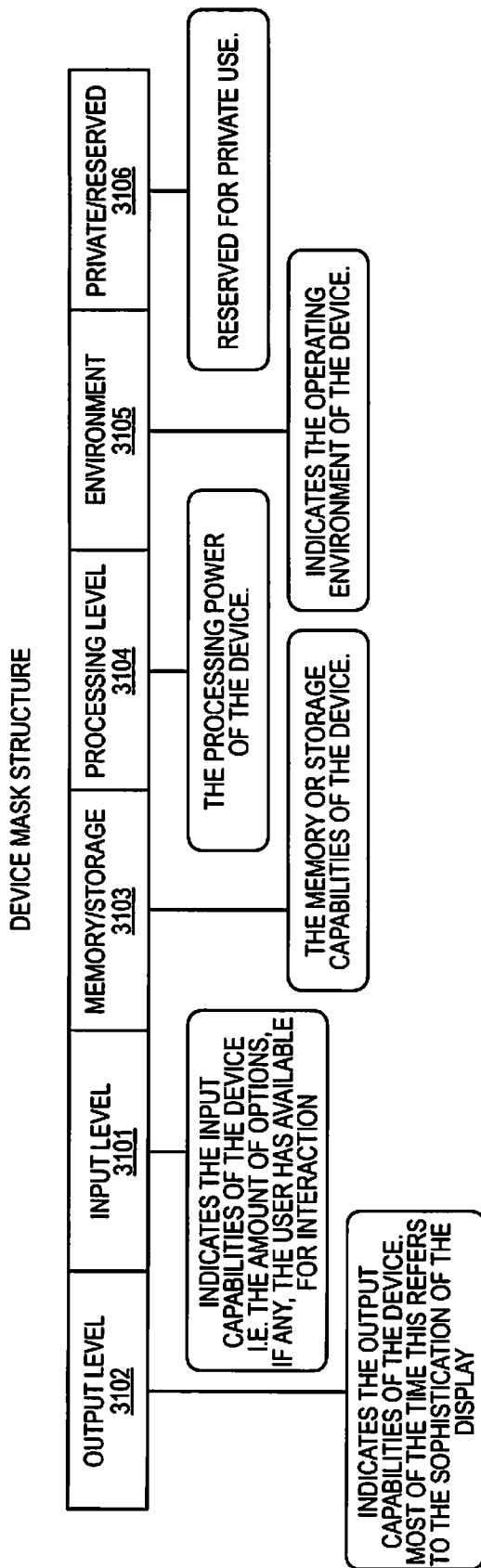


FIG. 31

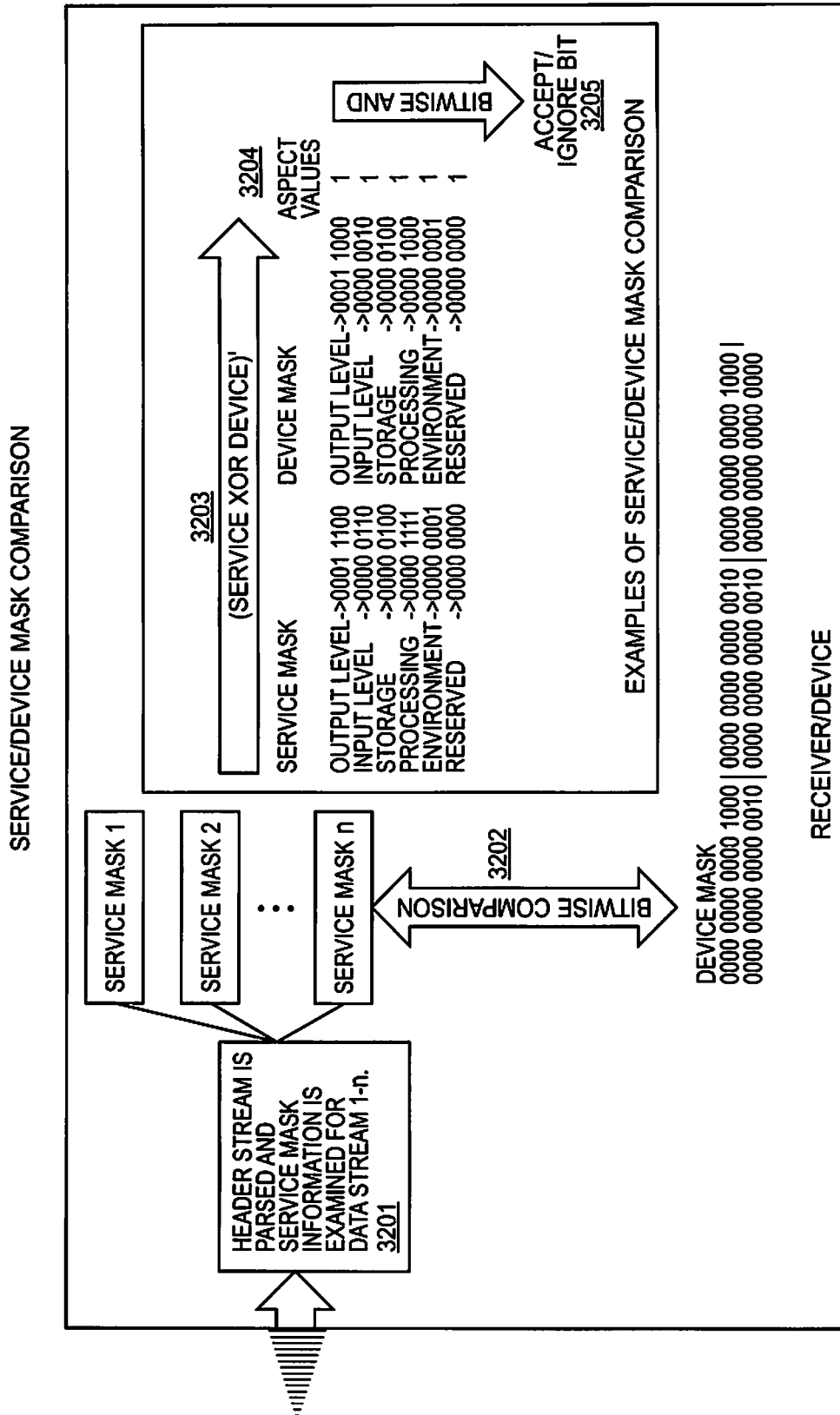


FIG. 32

SYSTEM AND METHOD FOR TRANSMITTING DIGITAL MULTIMEDIA DATA WITH ANALOG BROADCAST DATA

PRIORITY APPLICATIONS

This application is a continuation of U.S. patent application Ser. No. 10/484,518 filed Dec. 2, 2004, now U.S. Pat. No. 8,396,100, which is hereby incorporated by reference in its entirety.

The '518 application claims priority under 35 U.S.C. §119 to U.S. Provisional Patent Application Ser. No. 60/306,080 entitled "EXTERNAL NETWORK INTERFACE" filed in the name of Paul Signorelli on Jul. 17, 2001; and further is a continuation-in-part of U.S. patent application Ser. No. 09/839,451 entitled "SYSTEM AND METHOD FOR GENERATING MULTIMEDIA ACCOMPANIMENTS TO BROADCAST DATA" filed in the name of David Corts et al. on Apr. 20, 2001, now U.S. Pat. No. 7,908,172, which is a continuation-in-part of U.S. patent application Ser. No. 09/802,469 filed on Mar. 9, 2001 which, in turn, claims priority to U.S. Provisional Patent Application Ser. No. 60/188,050 filed on Mar. 9, 2000; this application is further related to U.S. Patent Application Ser. No. 60/346,785 entitled "SYSTEM AND METHOD FOR ASSEMBLING SUPPLEMENTAL DIGITAL DATA TO BE BROADCAST ON AN SIDEBAND OF AN ANALOG BROADCAST" filed in the name of David Corts et al. on Jan. 7, 2002; and U.S. Patent Application Ser. No. 60/346,784 entitled "SYSTEM AND APPARATUS FOR TRANSMITTING DIGITAL MULTIMEDIA DATA WITH ANALOG BROADCAST DATA" filed in the name of David Corts et al. on Jan. 7, 2002, the entirety of each of these applications being incorporated herein by reference.

FIELD OF THE INVENTION

The present invention is directed generally to multiplex communications, and more particularly to communicating messages over free space (i.e. a radio frequency (RF) sideband/sub-carrier or frequency mask) for reception at multiple destinations.

BACKGROUND OF THE INVENTION

In-Band On-Channel (IBOC) broadcasting is an emerging Digital Audio Broadcasting (DAB) technology, developed by IBIQUITY DIGITAL, INC., that enables existing radio broadcasters to transmit digital data over current analog transmission frequencies. Such radio broadcasters commonly employ amplitude-modulated (AM) and frequency-modulated (FM) bands for the transmission of audio broadcast data. IBOC technology has the ability to create a "hybrid" signal that can simultaneously send both the analog and digital data. U.S. Pat. No. 5,757,854, incorporated in its entirety herein by reference, discusses these capabilities in greater detail.

Digital data may be digitally-formatted data or digitally compressed analog data. Digital data may include processing instructions for rendering visual and/or audio components on, for example, an IBOC receiver. Such processing instructions may be used to render synchronized visual components, such as text and images describing artist or song title information for currently-broadcast songs on the analog band, news headlines, traffic reports or other information that would be of interest to a radio listener. The digital data may include audio components for presenting selectable audio data.

In an IBOC network, IBOC receivers recognize analog and digital data broadcast by IBOC transmitters, and present such

data to a user through a display and/or an audio output. The user may interact with the data and provide a response via the IBOC receiver to either a party operating the IBOC transmitter or a third party. Additional examples of digital data and its uses are described in the previously incorporated '172 patent.

In order to accommodate these various IBOC network functionalities, a protocol for the assembly, transmission and synchronization of such digital data is described.

SUMMARY OF THE INVENTION

The present invention relates to the data formats used to transmit digital data over traditional analog bands and other features enabled by IBOC technology.

One aspect of the present invention relates to the transmission of digital data, such as digital audio data, over predefined channels, such as AM or FM channels using known radio broadcast equipment.

Another aspect of the present invention relates to the successful transmission of digital data over analog bands using various synchronization protocols between the sender and the receiver.

Still another aspect of the present invention relates to providing sufficient security for the digital transmission so as to prevent the tampering or corruption of data by an outside source. The security process involves various encryption protocols and authentication procedures. Yet another aspect of the present invention relates to the transmission of a response from an IBOC receiver to an appropriate operation handler. Such operation handler may be a native handler, wherein an embedded module or procedure exists to service the request. In another embodiment the appropriate operation handler may be a non-native handler, wherein the service request is transmitted to another device for handling.

Still another aspect of the present invention relates to the creation of a Quality-of-Service (QOS) system, wherein a group of RF carrier bands is created around each central frequency available for broadcast. Since the reliability of data transmission decreases with RF carriers further from the central frequency, these RF carriers may be grouped according to the volume of data that can be successfully accommodated within a predetermined time. For example, digital data corresponding to a real-time sporting event (which may require continuous updates of digital data) may be transmitted over a more-reliable, high-volume RF carrier or set of RF carriers, while digital data corresponding to a weather report that is updated only every hour may be repeatedly transmitted over a less-reliable, low-volume or set of RF carriers to insure reception of all required digital data.

BRIEF DESCRIPTION OF THE DRAWINGS

Further aspects of the instant invention will be more readily appreciated upon review of the detailed description of the preferred embodiments included below when taken in conjunction with the accompanying drawings, of which:

FIG. 1 is an of an exemplary IBOC network including transmitters, receivers and third party service providers;

FIG. 2 is a diagram depicting an exemplary flow of information between the devices shown in FIG. 1;

FIG. 3 is a block diagram of an exemplary data format used in the transmission of digital data over the IBOC network of FIG. 1;

FIG. 4 is a block diagram of an exemplary organization data for the format of FIG. 3;

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FIG. 5 is a flowchart depicting an exemplary process for invoking a user agent handler to accomplish the reception of digital data;

FIG. 6 is a flowchart depicting an exemplary process for operating the user agent handler of FIG. 5;

FIG. 7-7B is a flowchart depicting an exemplary process for providing a request over the IBOC network;

FIG. 8 is a flowchart depicting an exemplary process for invoking the service handler to accomplish a request;

FIG. 9 is a flowchart depicting an exemplary process for validating a message transmitted through the IBOC network;

FIG. 10 is a block diagram of exemplary hardware and software used for message authentication with a one-way hash function;

FIG. 11 is a flowchart depicting an exemplary process for authenticating data received over the IBOC network;

FIG. 12 is a block diagram depicting an exemplary data structure for an authentication header;

FIG. 13 is a block diagram depicting an exemplary data structure for accommodating synchronization of digital and analog data;

FIG. 14 is a flowchart depicting exemplary operation of a static Quality-of-Service manager;

FIG. 15 is a flowchart depicting exemplary operation of an active Quality-of-Service manager;

FIG. 16 is a depiction of an exemplary JAVASCRIPT module for the Quality-of-Service manager;

FIG. 17A-B are depictions of an exemplary RF carrier javascript module;

FIG. 18 is a depiction of an exemplary RF carrier Factory javascript module;

FIG. 19A-B are depictions of an exemplary RF carrier Pool javascript module;

FIG. 20A-B are depictions of an exemplary Sub-channel javascript module;

FIG. 21 is a depiction of an exemplary Sub-channelFactory javascript module;

FIG. 22A-B are depictions of an exemplary Service javascript module;

FIG. 23 is a description of the ServiceListenerjavascript module;

FIG. 24A-C are descriptions of the ServiceMetadatajavascript module;

FIG. 25 is a flowchart depicting an exemplary process for using a synchronization cue to create an audio cul-de-sac; and

FIG. 26 is a flowchart depicting an exemplary process for using a synchronization cue to trigger an audio cul-de-sac.

FIG. 27 is an illustration depicting an exemplary channel header body and service mask;

FIG. 28 is an illustration depicting an exemplary channel header body and a Quality of Service filter mask;

FIG. 29 is an illustration depicting an exemplary processing of filter masks;

FIG. 30 further illustrates one non-limiting example of the practical effects of employing service masks;

FIG. 31 illustrates another non-limiting example embodiment of a device mask; and

FIG. 32 illustrates one non-limiting example of the comparison of service and device masks.

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DETAILED DESCRIPTIONS OF THE INVENTION

Broadcast Data

In one non-limiting example embodiment of broadcast data, a given radio frequency may carry the following information: a streamed analog broadcast, and an analog sub-carrier data, and/or the like.

Sub-carrier data is generally small text or numeric information. The sub-carrier data is carried on a standardized set of RF frequencies using any number of standard transmission technologies.

An RF carrier is a single carrier frequency on an AM or FM radio channel capable of carrying n-bits of data. For example, one RF carrier may carry 4 bits or 8 bits of data. RF carriers have varying degrees of robustness. The robustness of an RF carrier can relate to one or more factors including distortion from adjacent radio frequencies and distortion from the analog carrier on the same channel. To overcome this, on RF carriers with lower robustness, data can be given a greater emphasis on error correction. This can include looping (re-transmitting) data, using forward error correction techniques, and interleaving the data. Applying more error correction data has the effect of reducing the bandwidth. The robustness of a given RF carrier for a particular broadcast facility can be reasonably estimated and predicted. Consequently, over a given period of time, the bandwidth contribution of a single RF carrier or a series of RF carriers for a broadcast facility can be calculated.

With the advent of IBOC a given radio frequency will continue to carry a streamed analog broadcast, but also has the ability to have one or more streamed digital audio broadcasts as well. At least one of the digital audio broadcasts is intended to be a digital duplication of the analog audio broadcast. On top of that, the IBOC system allows for the transmission of binary and ASCII files, and the streaming of text and numeric information with the digital audio. Analog sub-carriers can co-exist with the new digital IBOC information on the RF carriers as well. Thus, at any given time a radio broadcaster could have a single streamed analog audio broadcast, one or more streamed digital audio broadcasts, a series of text and numeric information streamed with digital audio, any number of binary and ASCII files, sub-carrier data, and/or the like. Any of these can be considered broadcast data.

Furthermore, the streamed text and numeric information can carry instructions and data to be rendered to the receiver. The binary and ASCII files can be multimedia data such as textual file formats (e.g., ASCII plain text, rich text, html, xml, and/or the like), audio file formats, graphic file formats, video file formats (e.g. MPEGs, MP3s, JPGs, GIFs, and/or the like), multimedia file formats, and/or the like. These files can also be mark-up or instructions to an application on the receiver.

Trigger Event

In one non-limiting example embodiment, a trigger event is an event that occurs on the receiver that causes an action to occur. The action that occurs can be any action available to the receiver. Examples of receiver actions are displaying data, playing an audio file, recording the main program audio stream, pulling data from the storage, and/or the like. If the receiver also has a two-way communication channel associated with it, then the action could initiate communication on this second channel. In one implementation, there are three types of trigger events: a broadcast event, a system event, and a user event. A broadcast trigger event occurs via the broadcast itself. A broadcast trigger event associates a receiver action with some element of broadcast data. For example, it

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could be an event associated with the audio stream (analog or digital) that can be used to synchronize program audio and data. It could be an event associated with a particular time kept by the broadcaster. This could be used to tell the receiver and/or user that it is the start of the 7 A.M. broadcast hour or that it is 7 A.M. according to the National Institute of Standards and Technology.

A system event occurs when a receiver condition is met. System events can be the time kept by the receiver, the location of the receiver, and a signal receiver from a separate communication channel. For example, a system event can occur when a receiver in a car with a navigation system receives a notification from the navigation system as to the location of the vehicle. In such an example, the receiver can request a regional ad query so that regional advertising would be scheduled with broadcast data.

A user event is initiated through the receiver by the user. A user event could be the pressing of a button on the receiver or even a voice command issued by the driver.

Trigger events operate similarly. When they occur, they engaged associated actions.

FIG. 1 provides an overview of an exemplary IBOC communication network 100. A service provider 110 sends formatted digital data suitable for transmission to IBOC transmitter 120. The IBOC-formatted data may also be transmitted over other types of networks, such as a co-axial network or a fiber optic network. The IBOC transmitter 120 may then transmit the digital data, as well as analog audio data, over a wireless network to a plurality of IBOC user devices 140. The IBOC user device 140 may be an analog/digital radio receiver. The IBOC user device 140 may be configured to receive IBOC formatted data in addition to other data formats, such as wireless access protocol (WAP) formats. As seen in FIG. 1, the IBOC user device 140 may possess an external service interface, which allows it to receive data from a third party external network service 130, that may or may not be configured to transmit IBOC formatted data.

FIG. 2 is a diagram of the operation of the IBOC system in one embodiment. The data provider 210 is analogous to the service provider 110 of FIG. 1, which assembles and distributes the digital data for IBOC transmission. The digital data is then sent over a dedicated gateway link 220 to a broadcaster 230. The gateway link 220 may be a wireless link, such as a cellular link, or a may be a non-wireless connection such as a fiber optic link. In the IBOC scheme, the broadcaster 230 may be an existing radio broadcaster. By using existing broadcasters, new hardware for accomplishing radio transmission need not be introduced during the implementation of the IBOC system. The IBOC formatted digital data is then transmitted wirelessly to the receiver device 270 using traditional frequencies in the AM or FM band.

Besides transmitting through radio frequencies over existing networks, IBOC formatted data may also be transmitted over an intermediate entity 250. The intermediate entity 250 may be a telematics network, the internet, a cellular network, or a number of other data networks. Data going through this intermediate network 250 is transmitted wirelessly to the IBOC user device 270.

FIG. 3 is an example of the organization of digital data into a format for IBOC transmission. Elements 310-340 represent block headers that are transmitted prior to the payload digital data 350, such as digital audio data to be rendered for a user of the user device 140. The first part of the block header is datatype filter 310, which provides information on the type of data in the payload, such as audio data, visual or textual data,

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such as song title or artist name, or any number of other digital communication that can occur during the operation of the system.

Filter Masks

To better understand the channel headers, it is useful to understand the power and flexibility of filter masks. A filter mask is information that may be embodied in a "tag" that is assigned to data from a data service. This tagged filter mask may be used by a receiving device to either accept or reject various forms of data. In one non-limiting example embodiment, the ability and manner in which the filter mask is employed, i.e., its capability, depends on the abilities of a particular radio receiver. The filter mask can be and/or mask any kind of data. For example, certain devices have limited capabilities, which would benefit from filter masks that filter out data types that are not understood by the device. As such, an example IBOC capable radio that had no video and/or text display abilities might have a filter mask that enabled it to ignore any video and/or text display information. One of the numerous advantages to such a filter mask is that the limited example device could simply ignore information that is tagged for video and/or text and not expend resources to interpret and/or process such data. Such a filter mask can also be used to manage subscription based information; i.e., subscribing customers would be allowed to view certain data while non-subscribers would not. Thus in one non-limiting example embodiment, a device would maintain tags representing its various abilities and/or subscriptions as a series of tags, which would act as a filter mask, i.e., filtering information. In such an example, broadcast data would include data in tagged format so that it may be discerned and acted upon appropriately. As such, when a device receives such tagged broadcast data, the device will compare the received tagged data to its own abilities as represented by its maintained tags. Received tagged data that matches the devices maintained ability tags will be processed appropriately by the device. In one example, XML based tagging structure may be employed for such filtering and/or tagging. For example, tagged data may be of the form:

```

1 <?xml version="1.0" standalone="yes"?>
  <formats>
    <textual>
      <textual_data>Song Lyrics</textual_data>
      <title>A Song Title</title>
      <author>John Doe</author>
      <publisher>Acme, Inc.</publisher>
      <pub_year>1999</pub_year>
      <subscription_information>FALSE</subscription_infor-
        mation>
    </textual>
    <audio>
      <digital_data>DATA</digital_data>
      <title>A Song Title</title>
      <author>John Doe</author>
      <publisher>Acme, Inc.</publisher>
      <pub_year>1999</pub_year>
      <subscription_information>TRUE</subscription_infor-
        mation>
    </audio>
    <video>
      <digital_data>DATA</digital_data>
      <title>A Song Title</title>
      <author>John Doe</author>
      <publisher>Acme, Inc.</publisher>
      <pub_year>1999</pub_year>
      <codec>MPEG4</codec>

```

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```
</video>
</formats>
```

An example device might have the following have its capabilities defined as such:

```
<?xml version="1.0" standalone="yes"?>
<formats>
<textual><capability>FALSE</capability>
</textual>
<audio>
<capability>TRUE</capability>
<subscription_information>TRUE</subscription_infor-
  mation>
</audio>
<video>
<capability>FALSE</capability>
</video>
</formats>
```

In such an example, as the only matching tag format types are "audio," only the audio "DATA" will be processed on the receiving device. In an alternative embodiment, a device may examine its own "capability" tags and process data based on its abilities. In such an embodiment, if the device receives data tagged with one of its known capabilities, such data will be processed. The capabilities may further define a specific kind of software requisite to process the data; i.e., a "codec" tag may be provided so that the receiving device must contain and/or obtain the requisite and/or specified codec (e.g., "MPEG4") for processing received data. In another alternative embodiment, a device may filter information with a subscription tag. In such an embodiment, if the broadcast data is tagged with a, for example, "subscription_information" being "TRUE" (and/or set to a password and/or some other security information, e.g., encryption key, etc.) and the receiving device had its subscription tag enabled with a matching "TRUE" and/or password, then the received data will be processed. Thus, with the above example, audio information tagged requiring a subscription would only be played on devices capable of producing audio and possessing an enabled subscription tag; i.e., a subscription tag on the device set to "TRUE" and/or the requisite password.

Such filtering allows a point-to-multipoint broadcasting medium to provide better targeting of information to receiving devices.

In an alternative filter mask embodiment, an application is defined. For example, an application in the IBOC world would be "program associated data" or "PAD." A code is assigned to a PAD. In this way, anytime an application is transmitted, the PAD code acts as the filter. Any receiver designed to interpret PAD and/or associated codes would know to look for that code and employ it as a filter.

FIGS. 27 through 32 illustrate one non-limiting example embodiment of filter masks. A service mask is provided for each service, and as such may occupy a variable number of bits in a channel header body 2702 of FIG. 27. Quality of service filtering may add to the filter mask, providing another dimension of filtering 2802 of FIG. 28. Further, FIG. 29 illustrates that filter masks are processed as part of reading header channel data. The received filter mask is compared to the device filter mask (block 2911). If there is no match (at block 2912) between the received filter mask and the device filter mask, then the data service is read (block 2913). If there is a match (at block 2912) or if the data service is read (block 2913), then iteration continues to the next service mask (block 2914). If there are no more service masks (block 2915), then there is no more header channel data to read (block 2916). If there are more service masks (block 2915), then header channel data iteration continues by going to the

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next service mask (block 2917) and again comparing the service mask with a device mask (block 2911).

FIG. 30 further illustrates one non-limiting example of the practical effects of employing service masks. Various data providers A, B, and C supply their content (small text services, large text services and interactive data services, respectively) to a (iDAB) server for preparation for broadcast at a radio station WXYZ. After a program signal is transmitted, which now includes content from content providers A, B, and C along with regular broadcast materials, various devices obtain the signal and use the broadcast filter mask to enable the displaying of content matching the device mask in a receiving device. In one non-limiting example embodiment, a device mask may be supplied as a software upgrade to a device. Examples in FIG. 30 show that a PDA might be capable of displaying content from data providers A, B, and C because a PDA has an improved display and input capabilities allowing for long textual and graphic display as well as interaction; while a low-end car receiver is capable of only displaying content from data provider A as it may only have a limited alphanumeric display capable of showing short messages; and a high-end car receiver is able to display content from both data providers A and B because it has an enhanced screen display allowing it to display long messages and graphics.

FIG. 31 illustrates another non-limiting example embodiment of a device mask. As is shown, input capabilities 3101, output capabilities 3102, memory capabilities 3103, processing level 3104, device type environment 3105, and various reserved 3106 parameters may define a device mask 3107. In one non-limiting example embodiment, the input capabilities 3108 indicate the types of interaction options available to a user (e.g., level 0 having no interaction abilities, level 1 having the abilities to navigate back and forth in received data, etc.). Similarly, the output capabilities 3109 indicate the types of output options are available data (e.g., level 0 having no display, level 1 having less than 33 characters for display, etc.). The memory capabilities 3110 indicate the amount of memory the device has for storing data (e.g., level 0 having 512 Kb or less, level 1 having 1 Mb or less, etc.). The processing level 3111 indicates the amount of computational resources that are available to the device (e.g., level 0 being a slow device, level 1 a medium speed processing device, etc.). The device type environment segment 3112 of the device mask may indicate the general environment in which the device operates (e.g., a level 0 device being handheld, a level 1 device being used in a car environment, etc.) As is illustrated, many device and performance characteristics may be accounted for by making additional components and/or tags that comprise the device mask. As such, a component of the device mask may be reserved for other purposes 3113, thus expanding upon and providing for future operational flexibility.

FIG. 32 illustrates one non-limiting example of the comparison of service and device masks. A header stream is extracted from a broadcast and parsed 3201. Upon parsing out the various service masks, the receiving device may engage in a bitwise comparison between the parsed service masks and the device mask within 3202. By employing a bitwise exclusive-or (XOR) function on both the service mask and the device mask 3203, aspect values 3204 are obtained and may be used to accept or ignore 3205 data associated with a given service mask. Of course, numerous forms of service and device masks may be used (e.g., the XML tag form as already discussed), and various types of comparisons may be employed as well.

Channel Header Structure

Moving back to FIG. 3, the second part of the block header is the recipient authentication **320**, which provides a means for the sender to allow only a subset of receivers to accept the data. For example, the data may be part of a subscription-only service. Using the recipient authentication **320**, the data sender may format the digital data so that only subscriber receivers will be able to receive and render it. The authentication process may also employ a number of protocols such as one-way authentication, where only the sender transmits an authentication message; two-way authentication, where the sender transmits an authentication message and the receiver transmits a reply authentication message; or three-way authentication, where the sender transmits an authentication message and the receiver transmits a reply authentication message and the sender further transmits a reply to the receiver's reply authentication message.

The third part of the block header is the data provider authentication **330**, which provides information that the recipient may use to confirm that the data that has been received from an expected broadcaster. The data provider authentication **330** may employ the same one-way, two-way, or three-way authentication process described previously for the recipient authentication **320**.

The fourth part of the block header is the encryption information **340**, which contains information about the encryption algorithm used to encode the digital data. The encryption information may contain a public key or a number of other data for a number of other encryption algorithms.

Finally, payload **350** contains the digital data that is to be transmitted to and rendered by the IBOC user device **140**. Payload **350** may contain a plurality of sub-headers that can be used in conjunction with the headers described in FIG. 3.

Sender time stamp **404** is a 32 bit word representing the time the data was sent from the broadcaster **230**. In this embodiment, the time may be represented in milliseconds. For higher accuracy, the time may be represented in tenths of milliseconds or even nanoseconds, and the size of the sender time stamp may be increased accordingly. The receiver time stamp **406** is a 32 bit word analogous to the sender time stamp **404**. It represents the time of the last time the data was rendered/executed by the receiver. Like the sender time stamp **404**, the receiver time stamp **406** is expressed in milliseconds, but may be configured to more precision as the size of the time stamp is increased to be more than 32 bits.

A problem present in this system that is not present in the radio systems in the prior art is that the receiver would have to be in synchronization with the transmitter to insure that there is no skipping or other anomalies in audio playback. A series of synchronization events are used to insure proper synchronization, these events are contiguous series of time that are correlated with the audio broadcast. For example, the digital data to be rendered could be a 30-second advertising commercial to be played during the first 50-second interval of a song. A synchronization cue is used to trigger this synchronization event.

The Event ID **424** is used to identify the correct synchronization cues for any particular digital data. Finally, User Data **432** is the executable data, such as multimedia data, to be rendered.

Beside these fields in the header body, there are other optional fields that may facilitate data transmission in the present invention. Such fields include synchronization cue field **402** which may be, for example, a sixteen-bit word that may be used to place the data on demand so that the data can be readily rendered at the appropriate time.

Another optional field is a domain identification field **408**, which may include a 4 byte long word that, in the application for digital radio, may be used to broadcast the call letters of the particular radio station, thereby identifying the source of the transmission. The call letters may be "WNBC," "WXRQ," "WNEW" or any combination of four alphanumeric characters. These letters may be encoded as digital information using any encoding scheme, such as the American Standard Code for Information Interchange (ASCII) standard.

The next block is content rating field **410**, which may be a 4 bit nibble that can be used to designate a rating on the program being broadcast. This feature allows the user to exercise a certain degree of listener discretion to avoid certain objectionable materials. There are 16 different ratings that can be designated on a particular transmission, for example, a "0001" nibble may be encoded as being intended for a general audience, whereas a "1111" nibble may be encoded as being intended for an adults only audience.

After content rating **410**, the next block of data is content category field **412**. Currently there is a variety of programming available on radio, such as music programming, broadcasting of sporting events, talk shows, interviews, and public addresses. A digital radio receiver has a display means such as a liquid-crystal display (LCD) that can present information about the category of the transmitted programming. That information is encoded within the 5 bytes of the content category field **412**. The first byte of the content category **412** may be the most generic (i.e. music) with later bytes of the content category field **412** being more specific (i.e. rock-n-roll). In another example in which a sporting event is being broadcast as analog data, the first byte of the content category field **412** may be encoded to be "00001000" to indicate sports, the most generic category. The second byte may be encoded to be "00000010" to indicate that the sporting event is a Major League Baseball game, a narrower category than the first byte. The third byte may be encoded to be "00000001" to indicate that the home team in the baseball game is the Yankees, a category narrower still. The fourth and fifth byte, in this example, would not be applied. The same structure can be applied to music. When music programming is being transmitted, the first byte of the content category field **412** may be encoded to be "00100010" to indicate music, the most generic category. The second byte may be encoded to be "00010001" to indicate rock music, a narrower category than the first byte. If the music being transmitted is classical music, then the third byte may be encoded to be "00001100" to indicate Baroque music, or other values to indicate Romantic, Medieval, or Modern music, a category narrower than the second byte. The fourth byte may be encoded to be "00110011" to indicate that the music being transmitted is a chamber piece, or other values to indicate an orchestral piece, a violin solo, or any number of other types of performances. Similar categories can be designated to other types of radio programming being transmitted, such as talk shows or public addresses.

File size number field **414** and file size magnitude field **418** are the next optional blocks in the header, following the content category field **412**. File size number field **414** may be a 16 bit word that indicates the size of the file being sent, for example, in kilobytes. File size magnitude field **418** may be a 2 bit word that indicates the magnitude of the file size number (i.e. bits, bytes, kilobytes, megabytes). The size of the file may be obtained by examining the file size number field **414** and the file size magnitude field **418**. For example, if a file being sent is 5 megabytes in size, the file size number **414** may be encoded as "0000000000000101," the number "5" in the

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binary system, and file magnitude **419** may be encoded as "11," which may represent megabytes in the current system.

Reserved bits field **420** are optionally allocated for future use as the digital radio system is dynamically designed so changes can be readily incorporated in the current system.

Status flags field **422** may be used during the operation of the digital radio to set flags for various data used by the user device **140**.

Group identifier (ID) field **428** can be used to identify stored digital data blocks that can be combined and re-used. This, in turn, decreases the data load to be transmitted by broadcaster **230** by reducing redundant data transmission. FIGS. **5-9** depict the operations of a user agent module employed by the user device **140**. FIG. **5** is a flowchart depicting a user agent handler invocation process **500**, the first process encountered by the incoming data at the receiver end. The user agent first parses the incoming digital data (step **510**). The parsing may be done according to frequency, data type, modulation (such as AM or FM), or any number of other useful parameters. The agent then searches the parsed data (step **520**). If the data relates to an external command, (step **550**), then the user agent invokes the message handler appropriate for the external command (step **570**). If the data relates to none of the external commands, then an invalid command message is produced (step **560**). After the determination is made on whether if the data relate to external command and appropriate actions taken at either step **560** or **570**, the invocation process returns to step **520** and iterates through the rest of the incoming data. During the iteration process, the user agent checks for the end of the data stream at step **530**. An indication for the end of the data stream may be an end-of-file (EOF) flag in status flag field **422**. When the end of the data stream is reached, the user agent ends the handler invocation process (step **540**).

FIG. **6** is a flowchart depicting the operation of the user agent handler after it has been invoked at step **570** above. The user agent handler first receives a notification message from the user agent (at step **610**). It then processes the information received at step **620** and builds an extensible mark-up language (XML) service invocation message corresponding to the incoming message (step **630**). It should be noted that although the invocation message is described to be written in XML in this embodiment, the message may also be written in HTML or other mark-up languages. The invocation message is then sent to the service register (step **640**).

The process **600** then continues in a series of iterative steps in which the handler waits for the response from a service register module, at steps **650-680**. If either the service register responds prior to the timeout of the handler's waiting period (step **660**), or if the handler's waiting period times out (step **680**), the process **600** ends (step **690**). The timeout threshold may be set statically, such as 45 seconds, in which responses from all service registers are required to be submitted to the handler in that time span. The timeout threshold may also be set dynamically, so that a different threshold is set depending on which service register is requested.

FIG. **7** is a flowchart depicting an exemplary process **700** performed by a native service register. A native service register must have embedded within it a executable module or procedure that is stored by the user device **140**. According to the process **700**, the user device **140** first receives the invocation message sent in step **640** of FIG. **6** above (step **710**). It then checks to see if the incoming message is a request for service (step **720**). If not, then the operation of service registers moves back to step **710** to wait for a next invocation message. If so, then the register executes the request at step **730**. If the execution of the request is successful, when

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checked at step **740**, a success message is built at step **750** and the result is returned at step **770** to the user agent handler. If the execution of the request is a failure when checked at step **740**, an error message is built at step **760** and the result is returned at step **770** to the user agent handler. After the result is sent to the user agent handler, then the process **700** ends at step **780**.

FIG. **7B** illustrates non-limiting alternative embodiment service invocation instead of request **730**, a request handler is invoked **730b**. Upon invoking the handler, the system determines if the handler exists **740b**, if the handler does not exist the request is ignored **760b** and illation continues **710b**. If the handler does exist **740b**, a request is sent to the handler **750b** and illation continues **710b**.

FIG. **8** is a flowchart depicting an exemplary process **800** performed by a non-native service register. A non-native service register may not have embedded within itself an executable module or procedure, but rather it may be able to outsource the service request to other modules, which may or may not include other native service registers or a communication with an external device over a wireless or hard-wired network connection. Process **800** begins at step **802** wherein the non-native service handler receives the invocation message sent by the user agent handler at step **640** of FIG. **6** above. It then checks to see if the incoming message is a request for service (step **804**). If not, then the operation of service registers moves back to step **802** to wait for the next invocation message. If so, then the request message is parsed according to the services required at step **806**. Each parsed message is then validated at **808**. If the message is determined to be valid at step **810**, it is used to invoke a service module at step **812**. The result of the service is monitored and stored at step **814**. A response message is written with the result information at step **816** and at step **820** the response message is sent to the user agent handler. If the message is determined to be invalid at step **810**, then an error response is built at step **818** and at step **820** the error response message is sent to the user agent handler. After the result is sent to the user agent handler, process **800** ends (step **820**).

FIG. **9** is a flowchart depicting a process **900** by which a request message is validated. This validation process **900** is called at step **808** in FIG. **8**. First, the request is sent by the service register by received by the validation module at step **910**. The request message then undergoes a series of tests to see if the message has a valid service reference, a valid parameter mask, and a valid request so that the service is executable by an available software module. These tests are performed at steps **920**, **930**, and **940**, respectively. If any of the test fails, the message is determined to be invalid at step **950**, and an invalid result is returned (step **970**). If the request message passes all the tests, the message is determined to be valid at step **960**, and a valid result is returned (step **970**). It should be stressed that the three tests in steps **920**, **930** and **940** are only exemplary tests. Other embodiments of the present system may contain other tests or have fewer tests without moving away from the spirit of the inventive system.

In one embodiment of the current system, there are features that require sufficient security to insure that there is no data tampering when communications are in transmit. For example, in an embodiment of the present invention where the user is able to purchase audio CD's using the digital radio system, communications regarding the purchase order to and from the user, such as the initial order request from the user and the order confirmation to the user, should be adequately encrypted and authenticated. The types of authentication used in the present system may be a one-way authentication, a two-way authentication, or a three-way authentication.

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In a one-way authentication, the sender transmits a timestamp, a nonce value, and a particular user's private key along with the payload data. A nonce value is a temporary value unique to all valid authenticated data. The receiver may then authenticate the information using the public key equivalent

In a two-way authentication, in addition to going through one-way authentication, after decrypting the transmitted data, the receiver transmits a reply message to the sender containing a new timestamp, the original nonce value, and a new nonce value. The reply message will be encrypted with the sender's public key encryption, which the sender may decrypt with the corresponding private key.

A three-way authentication may be used if the sending device and the user device 140 have not achieved synchronization. In addition to going through two-way authentication, after receiving the reply message, the sender transmits another reply to the receiver, containing the nonce value included in the first reply. After matching nonce values, the user device 140 may disregard the timestamps.

FIG. 10 is a block diagram depicting an exemplary one-way authentication process 1000 using a hash function. At step 1010, a message is first processed with a hash function to create a message digest 1020. The message digest is then encrypted using a private key K_1 (step 1030). The private key K_1 is then added to the encrypted message as a header and a message digest with header is produced (Step 1040). The message digest with header is then attached to the original message to produce the message packet at step 1050, which is then transmitted to the user device 140 (step 1060). The message digest with header and the original message are then extracted from the received message packet by the user device 140. The header of the message digest will be detected by the corresponding public key K_2 (step 1070). The detected header will be used to re-hash the received original message (step 1080), and the re-hashed message digest is compared with the received message digest (step 1090). If the message is determined to be authentic then the re-hashed message digest will be the same as the received message digest. If the two do not match, then the message is determined to be inauthentic and the message is discarded. This effectively prevents tampering by an outside party during the transmission of the message, since changing of even one bit of the original message may result in a significantly different hash which will not be recognized by the user device 140.

FIG. 11 is a flowchart depicting an exemplary process 1100 for authentication performed by the user device 140. The user device 140 first detects an incoming signal and reads it to determine the current authentication mode (step 1102). Next, the user device 140 detects and reads the timestamp value, the nonce, and various other header information necessary for authentication and stores them in memory (step 1104). The receiver then may confirm the authentication mode (step 1106). If the current authentication mode is the one-way mode, the receiver then proceeds to step 1114, discussed further below. If the current authentication mode is not one-way, then the authentication mode may be two-way or more, such as three-way. In such event, the process 1100 continues to step 1108 where the user device 140 transmits to the sender a new timestamp, the original nonce and a new nonce, as per the process of two-way authentication discussed previously above. The receiver then waits for a reply message from the sender (step 1110). The reply message may be an acknowledgment from the sender for receiving the new timestamp and the new nonce, or it may be an additional nonce for further multiple-way authentication. The receiver makes that differentiation at step 1112.

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If the reply is an acknowledgement, then the authentication mode is determined to be two-way authentication, and the receiver proceeds to step 1114. If the reply is not an acknowledgement, the user device 140 examines the incoming reply to determine whether the current authentication mode is three-way at step 1116. If the reply is not further authentication information, then the process moves once again to step 1114. If the reply contains further authentication information, then the receivers transmit additional nonces (step 1118) and will receive additional replies by moving back to step 1110 after step 1118. It is noted that although in the current discussions the most numerous multiple-way authentication mode is three-way mode, n-way authentication mode can be achieved by executing the iteration comprised of steps 1110, 1112, 1116 and 1118 n times.

At step 1114, the receiver detects and reads the hash algorithm ID, the key length, any digital signature information, and any other data necessary to decrypt the incoming message. These data fields will be discussed at length in the description of FIG. 12. If the authentication mode of the incoming message is two-way or more, then the timestamp information and the nonce information from the sender and the receiver will be compared as the first step of authentication (step 1120). Otherwise, the message is discarded if the timestamp and nonce information do not match. If the information does match, then the incoming un-hashed message is hashed at step 1122 and the incoming hash digest in decrypted at step 1124. The two results are compared at step 1126. If the results are equal, then the data is authentic and the message is passed onto the various other parts of the user device 140, such as the user handler, at step 1128. If the results are not equal, then the data is determined to be inauthentic and the message is discarded at step 1130.

FIG. 12 is a block diagram depicting an exemplary authentication header 1200 that may be used in the process 1100 depicted in FIG. 11. The different fields 1202-1238 in the header 1200 may be arranged in the same order that they will be detected by the receiver.

The authentication mode field 1202 is used to indicate the current authentication mode, that is whether it is one-way, two-way, three-way, or other multiple-way modes.

The timestamp 1206, first nonce value 1210, and second nonce value 1214 are information used in the authentication process 1100. Although in this particular embodiment, two nonce values are allotted, it would be fairly clear to one of ordinary skill in the art to increase the number of nonce value fields in the header so to increase the multiple-way authentication that may be used, as described previously herein.

Hash algorithm ID 1218 is a code for the hash algorithm used by the sender. This is used by the user device 140 to hash the un-hashed message in order to compare it against the decrypted hash digest message. The key length 1224 indicates the length of the public key that will be used to decrypt the hash digest, while public key 1228 contains the actual public key.

Digital signature length 1232 indicates the length of the digital signature and the digital signature field 1236 contains an indication of the user's identity which may be encrypted or signed by the sender. Items 1204, 1208, 1212, 1216, 1220, 1226, 1230, 1234, 1238 are examples of each of respective fields 1202, 1206, 1210, 1214, 1218, 1224, 1228, 1232, and 1236 in the authentication header.

Quality of Service (QoS) Management

Different degrees of reliability with respect to transmission of data over an IBOC broadcast are termed "Quality-of-Service" (QoS). In an IBOC system, a data channel is composed of an infinite number of RF carriers, that are specific frequen-

cies adjacent to central frequency over which analog data is broadcast. A finite number of these RF carriers may be used to reliably broadcast over a reasonable distance. Mainly due to the susceptibility of interference, some RF carriers can send large amounts of data over long distances with little chance of error, while others cannot. The RF carriers become more unreliable depending upon their proximity to central frequency analog data, or adjacent frequency data.

Since the entirety of the data channel is available for broadcasting, the broadcaster may subdivide the data being broadcast into data that needs to be transmitted with high speed and efficiency and that which does not require high speed or efficiency.

The focus of QoS system is the management of these sub-channels, and corresponding pricing for data, such as advertisements, based on the reliability and speed of different parts of the channels. Data services will request sub-channels for broadcast of equal sized data service packets containing data blocks. A data service is a collection of similarly purposed data, such as a communication of data between two parties. A data service packet represents a single unit of data for a particular data service and assembled from the individual data service packet segments, being all of the data within the data block for a particular service. A data service packet transmission may be interleaved over time to reduce distortion of the data being broadcast. A data block is a physical series of data bits, created from one or more of the RF carriers on a radio channel over a given period of time. The beginning of the block may be indicated by a recognizable signal pattern, which may be referred to as a synchronization pulse. The data block will consist of all the radio frequency carriers being read by the user device **130** for n^{th} millisecond over a larger period of time. All of the data for a block for a given RF carrier may be referred to as a block segment. A single data block may carry data for more than one data service, which in turn may be identified by a series of header bits.

The reliability of a sub-channel can be determined by the number of RF carriers used, or the size of the sub-channel, and the positions of the individual RF carriers on the spectrum in relation to the main data channel. Each RF carrier may be indexed 1 through n . The most robust RF carrier is 1, and a formula is used to determine the relative reliability of each subsequent RF carrier. The result of this formula is referred to herein as a "QoS Rating". In an exemplary method, the average QoS Rating of the sub-channels may be used to determine the "QoS Level" of that sub-channel. Given both the desired number of sub-channels and the desired QoS Level for each sub-channel, as requested by the data service, such sub-channels may be dynamically allocated in order to satisfy that demand.

A transmission proxy acts as the controller between data services and the data channel. A data service will make a request to the transmission proxy with all of the parameters for sending the data, encapsulated in a data service object. The transmission proxy then communicates with the QoS Manager (described below) for a sub-channel object, and thereby inserts the data into the channel.

In an exemplary process, there exists three methods by which RF carriers can be grouped into sub-channels. These are as follows: (i) static, wherein the sub-channels are pre-defined before a data service is given access; (ii) dynamic, wherein the sub-channels are defined at the time of the request for a QoS Level sub-channel; and (iii) hybrid, wherein the selection of RF carriers is both static and/or dynamic. In a static scenario, the data service makes a request for the sub-channel with a desired QoS Level, and if that sub-channel is

available it is returned. Otherwise, an error message is returned. If a valid handle is returned, the data service will then use the sub-channel to transmit data, and the sub-channel is considered allocated. Upon completion of the data service use of a sub-channel, the sub-channel does not have to be freed because it is statically defined.

In a dynamic scenario, if a collection of RF carriers is available produce a sub-channel with the desired QoS Level at the time of the request, the RF carriers are grouped and the handle for the sub-channel is returned. Otherwise, an error message is returned. When the data service has completed using the sub-channel, it must be returned in this dynamic environment. In a hybrid scenario a pre-defined number of sub-channels are static and the remaining RF carriers utilize a dynamic scenario.

The QoS Manager interacts with a transport system, or other communication system such as a modem operating in conjunction with an application program interface (API), in order to gain access to sub-channels and allocate them appropriately. The transport, or modem may be any device that creates the waveforms on the RF carriers. The API provides a low level interface to the functions of the modem, such as asking for percentages of a given sub-channel, or acknowledging the status of a given sub-channel. The QoS Manager must be able to create all of the three above mentioned groupings of RF carriers into sub-channels and maintain a pool of RF carrier objects produced by the modem that hold information regarding the reliability and status of the RF carrier or sub-channel it represents.

There exist at-least three statuses for a sub-channel object: busy; available; and killed. A static object will be killed when the dedicated configuration has been compromised by other dedicated configurations. A dynamic object will be killed when the channel can no longer support that level of bandwidth without compromising service on other sub-channels. Eventually these killed sub-channels will be removed from the pool of sub-channels by the QoS Manager. The object may also have a lease time, during which it becomes busy and no other process may use it. If the object has not been returned to the pool by the end of the aforesaid lease time it will automatically be returned to the pool by the QoS Manager. The QoS manager may also periodically recycle sub-channel objects when they are not in use, which requires destroying and re-building them.

FIGS. **14** and **15** describe the exemplary structure and method for initialization, selection and assignment of sub-channels and RF carriers based on Quality of Service definitions according to the present invention. FIG. **14** displays an exemplary process **1400** for initialization of sub-channels based on static parameter definitions which are examined at step **1402**. From these definitions, the available sub-channels are determined at step **1404**. Furthermore, the sub-channels are stored in an addressable memory structure at step **1406**, after which process **1400** ends.

FIG. **15** provides an exemplary process **1500** for the selection and assignment of sub-channels based on parameters defining the quality of service desired, which are examined at step **1508**. These parameters are used to determine whether the request for a sub-channel is static or dynamic (step **1510**). If the request is static, the availability of that specific channel is further determined at step **1512**. If the requested channel is available, it is retrieved from an addressable memory structure at step **1514**. The addressable memory structure used at step **1514** may be the same memory used at step **1406**. Following retrieval, a lease time is set and the sub-channel identification is returned to the entity that requested it at step **1516**.

If, on the other hand, the requested channel is not available, an error message is returned to the entity that requested it at step **1518**.

If the request is for a dynamic sub-channel at step **1510**, the available sub-channels are determined at step **1520**. The ability to create the requested sub-channel is further determined at step **1522**. If the requested sub-channel can be built then a lease time is set and the sub-channel is returned to the entity that requested it at step **1516**. If the requested sub-channel cannot be built, an error message is returned to the entity that requested it at step **1524**.

FIG. **16** is a description depicting an exemplary Java interface structure for the QoSManager **1610**, that handles the creation of sub-channels and the level of service for a sub-channel. Furthermore, it is the point of entry for data services requiring a sub-channel to send data. The interface QoS Manager includes a getSub-channel() command **1620** that returns a sub-channel of a specific QoS level, or returns null if the specified sub-channel cannot be returned. In addition, the method detail of the QoS Manager **1630** provides exemplary parameters for each element in the interface device method summary.

FIGS. **17A-17B** are descriptions depicting an exemplary Java interface structure for an Interface RF carrier **1702**, that handles the reading and writing of information onto a particular RF carrier comprising a groups of sub-channels. The interface RF carrier **1702** first involves clear() function **1704**, which is an initialization function that empties the RF carrier of all data and returns it to the RF carrier pool. The getQoS-Rating() function **1706** assigns a rating to the RF carrier based on reliability. This rating may be used to determine the speed of data transfer onto the RF carrier as well as the data volume. The read() function **1708** reads data from the RF carrier, while the write() function **1710** writes data onto the RF carrier.

The getQoSRating() **1712** function is a public function that has no input arguments and may return the RF carrier's rating as an integer. The read() function **1714** is a public function that has argument array which includes the data read from the RF carrier and an integer representative of a number of bytes requested to be read. It returns the number of bytes actually read as an integer.

The write() function **1716** is a public function that has argument array which is the data to be written to the RF carrier. It returns the number of bytes actually written as an integer.

The clear() function **1718** is a public function that has no input arguments and returns nothing. Instead, it initializes the RF carrier by emptying it of all data and returning it to the RF carrier pool so that it can be available for future use.

FIG. **18** is a description depicting an exemplary Java interface structure for an Interface RF carrierFactory **1802**, which handles requests to create new RF carriers within the transmission band. The interface RF carrierFactory **1802** involves the function newRF carrier() **1804** which creates a new RF carrier with a specified QoS rating.

The function newRF carrier() **1806** is a public function that has an integer input argument rating, which specifies the QoS rating of the new RF carrier. This function returns the object RF carrier, which may be a pointer or a memory location with floating data type containing the frequency of the new RF carrier.

FIGS. **19A-19B** are descriptions depicting an exemplary Java interface structure for the Interface RF carrierPool module **1902**, which handles the locking and unlocking of RF carriers for use by a sub-channel. This interface prevents

unnecessary reconstruction of RF carriers by the interface RF carrier Factory **1802** previously discussed.

The interface RF carrier Pool **1902** involves the checkIn() function **1904**, which is a function that notifies the RF carrier pool that a particular RF carrier object passed through as its input argument is available for checkout.

The checkout() function **1906** moves a RF carrier with a particular QoS rating from the RF carrier pool to use by a sub-channel. The GetAvailableRating() function **1908** returns the QoS ratings of the available RF carriers currently in the pool. The function GetCount() **1910** returns the current size of the RF carrier pool. The method detail of the interface RF carrier Pool **1902** further describes the functions used by the interface. The function checkout() **1912** is a public function that has an integer input argument qosRating, which specifies the QoS rating of the requested RF carrier. A RF carrier object is returned. The function throws to an error message when no RF carrier with the input QoS rating is available in the pool. The function checkIn() **1914** is a public function that has a RF carrier object input argument subc whose frequency, QoS rating, and other attributes will be listed in the pool so that the RF carrier is made available for checkout. The function returns to an error message if there is already a RF carrier in the pool with the same QoS rating, so to reduce unnecessary check-ins and check-outs.

The function getAvailableRatings() **1916** is a public function that has no input arguments, but it returns an integer array storing the available QoS rating when the function is called.

The function getCount() **1918** is a public function that also has no input arguments, but returns an integer that represents the current size of the RF carrier pool, or the number of RF carriers available, when the function is called.

FIGS. **20A-20B** are descriptions depicting an exemplary Java interface structure for the interface sub-channel **2002** which handles data services to send data to a channel or receiving data from a channel. The interface sub-channel **2002** includes a destroy() function **2004**, which removes all data from a particular channel and returns all the RF carriers that the sub-channel is comprised of.

The function getInputStream() **2006** gets the data stream that is used to read data from the sub-channel. The function getOutputStream() **2008** gets the data stream that is used to send data to the sub-channel. The function getQoSLevel() **2010** gets the QoS rating of the sub-channel currently in use.

The function GetOutputStream() **2012** is a public function that has no input arguments. When called, it returns an OutputStream object that contains the data that will be sent to the sub-channel. The function returns an error message if the output stream cannot be returned. Such a condition occurs if the output stream does not exist or is otherwise busy.

The function GetInputStream() **2014** is a public function that operates similarly to the function GetOutputStream() **2012**. GetInputStream() **2014** also has no input arguments. When called, the function returns an InputStream object that will be used to store the data read from the sub-channel. The function returns an error message if the InputStream object cannot be returned. Such a condition occurs if the input stream does not exist or is otherwise busy.

The function getQoSLevel() **2016** is a public function that has no input arguments and returns the QoS rating of the sub-channel as an integer.

The function destroy() **2018** performs the cleanup work required to inactivate a sub-channel. When called, the function removes all data from the sub-channel, disassembles the sub-channel into various RF carriers and returns an array of integers that indicate the particular RF carriers that were extracted from the sub-channel. These RF carriers may be

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made available to other sub-channels once they are returned to the RF carrier pool using such interface as the RF carrier-Pool interface **1902**.

FIG. **21** is a description depicting an exemplary Java interface structure for the interface Sub-channelFactory **2102**, which is called to create new sub-channels from available RF carriers. The interface uses the function newsub-channel() **2104** to create a new sub-channel using a particular group of RF carriers. The function newSub-channel() **2106** is a public function that has an input argument of an array of RF carrier objects. These RF carrier objects are to be used by the function to create the new sub-channel. This function **2106** returns a sub-channel object that is ready to be used in the system by such interfaces as interface SubChannel **2002**. The function **2-16** generates an error message when the sub-channel cannot be created by the RF carriers indicated by the input array. This may occur when the RF carriers indicated by the input array is not available in the RF carrier pool, such as when these RF carriers are used by another sub-channel.

FIGS. **22A-22B** are descriptions depicting an exemplary Java structure for the interface Service **2202**, which allows a receiver to handle the incoming information. The interface Service **2202** involves the function Authenticate() **2204**, which confirms the authenticity of the sender by matching timestamps and nonces with the process described by FIG. **11**. The function getInputstream() **2206** is the same process as the function **2206** described previously. Similarly, the function getOutputStream() **2208** is the same process as the function **2008** described previously. The function getService-MetaData() **2210** is used to read the incoming header and other information. The specifics of the function will be discussed in the descriptions for FIGS. **24A-24C** below. The function setServiceMetaData() **2212** reads the information received by getServiceMetaData() **2210** and configures the receiver unit accordingly. The function getServiceMeta Data() **2214** is a public function that has no input arguments and it returns a class of type ServiceMetaData when called. The particular fields in the class ServiceMetaData will be discussed in the descriptions for FIGS. **24A-24C**. The function setServiceMetaData() **2216** is a public function that has no input arguments and returns nothing. When called, it configures the receiver according to the ServiceMetaData class recorded by getServiceMetaData() **2214**. The functions getOutputStream() **2218** and getInputStream() **2220** are the same functions as getOutputStream() **2012** and getInputStream() **2014**, described previously above. The function Authenticate() **2222** is a public function that has the input argument of type "devicekey" which is an object that maybe used identify the particular sender of the transmitted data. The function returns nothing, but it throws to an error message if unauthentic information is found.

FIG. **23** is a description depicting an exemplary Java structure for the interface ServiceListener **2302**, which is responsible for delegating the construction of sub-channels and handing them off to the appropriate handler objects. To accomplish this, the interface ServiceListener **2302** must have access to interfaces RF carrierPoll **1902** and Sub-channelFactory **2102**. The interface ServiceListener **2302** also requires data on the services currently that are actively being received. Such data is supplied by the function getServices() **2304** that returns an array that contains information on the services that are actively being received.

FIGS. **24A-24C** are descriptions depicting an exemplary Java structure for the Interface ServiceMetaData **2402**, which extracts the header information from the incoming data stream described in FIG. **4**. The interface ServiceMetaData **2402** uses a separate function to extract each block of header

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described in FIG. **4**. The function getCategory() **2404** is a function that extracts the Content Category **412** in the header. Because Content Category is a group of 5 bytes indicating 5 levels of content scope, getCategory() **2404** has an input argument of integer type so that only the byte associated with the specified level is returned. The function getContentRating() **2406** is a function that extracts the Content Rating **410** in the header. The function getDataSize() **2408** is a function that extracts the File Size Number **414** in the header. The function getDataSizeMagnitude() **2410** is a function that extracts the File Size Magnitude **418** in the header. The function getDomainID() **2412** is a function that extracts the Domain ID **408** in the header. The function geteventIndicator() **2414** is a function that extracts the Event Indicator **426** in the header. The function getGroupID() **2416** is a function that extracts the Group ID **428** in the header. The function getMimeType() **2418** is a function that extracts the Mime Type **430** in the header. The function getReceiverTimeStamp() **2420** is a function that extracts the Receiver Time Stamp **406** in the header. The function getReserved Bits() **2422** is a function that extracts the reserved bits **420** in the header. The function getSenderTimeStamp() **2424** is a function that extracts the Sender Timestamp **404** in the header. The function getStatusBits() **2426** is a function that extracts the Status Bits **422** in the header. The function getSyncCue() **2428** is a function that extracts the Synchronization Cue **402** in the header. The function setReceiverTimeStamp() **2430** does not extract any information from the header, but configures the receiver so that the receiver will transmit the timestamp stored in the input argument "tStamp." After all the information is extracted from the incoming data, the interface ServiceMetaData **2402** stores the extracted data in a class with fields that correspond to the functions used by the interface, and that each field stores the header data extracted from the incoming data stream by the corresponding functions. This class can be used by other interfaces, such as interface service **2202**.

In the operation of radio systems including the present invention, there are numerous repetitions of the same audio data. Music is often repeated numerous times in the course of one day, and commercials can be repeated numerous times in a matter of minutes or hours. Therefore, it would be prudent for the receiver to record a number of these audio data in temporary memory location so to reduce the transmission data load while maintaining the operation of the system. The present invention accomplishes this data recording by transmitting a synchronization cue to create an "audio cul-de-sac" that is recognized by a receiver, such as user device **140**, and stored in a buffer or memory of the same. The audio cul-de-sac stored by user device **140** may also be multimedia information that a user can recall at any desired time.

FIG. **25** is a flowchart depicting the operation for using a synchronization cue to create an audio cul-de-sac. It is advantageous to use a synchronization cue to create an audio cul-de-sac because the synchronization cue contains data that identifies the multimedia data that will be recorded. The receiver first waits for a synchronization cue at step **2502**. When a cue is received, a check is performed at step **2504** to ascertain if the cue is a cue to signal the start of recording. If not, the user device **140** returns to step **2502**. If so, then the user device **140** starts to record multimedia data at step **2506**. While the multimedia data is being recorded, the user device **140** may also check for another synchronization cue at step **2508**. If the cue is determined to be an end cue at step **2510**, then the user device **140** stops recording the audio data. If the cue is not an end cue, then buffer check is done by the receiver at step **2512**. If the buffer is determined to be full at step **2514**, then the receiver also stops recording. If the buffer is not full,

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then the receiver returns to the state at step 2508 and waits for a synchronization cue. This process insures that the recording of multimedia data is stopped when an end cue is received or when the buffer has reached full capacity, which ever occurs first. After the recording, the data file in the buffer is marked with the synchronization cue that started the recording process, at step 2516, so that the file is identified, and the file is saved in cache at step 2518.

After the audio cul-de-sac is recorded, synchronization cues can also be used to trigger the cul-de-sac so that the audio file is played. FIG. 26 is a flowchart depicting the operation for using a synchronization cue to trigger an audio cul-de-sac. The user device 140 first begins by waiting for a synchronization cue at step 2610. After a cue is received, a determination is made at step 2620 as to ascertain if the received cue is a playing cue. If not, then the user device 140 returns to waiting for a cue at step 2610. If so, then the cue is read and the file ID contained within the cue is determined at step 2630. A search is performed at step 2640 to find a recorded audio file that has a file ID matching the one found in the playing cue. If an audio file is found at step 2650, then the file is played at step 2660. If not, then the receiver returns to waiting for a cue at step 2610.

In one non-limiting example embodiment, text and other media may be substituted for the audio data in the audio cul-de-sac. For example, a broadcaster may use a synchronization queue to synchronize piece of text to broadcast.

Although the invention has been described in detail in the foregoing embodiments, it is to be understood that the above descriptions have been provided for purposes of illustration only and that other variations both in form and detail can be made thereupon by those skilled in the art without departing from the spirit and scope of the invention, which is defined solely by the appended claims.

What is claimed is:

1. A method of transmitting digital data on at least one RF carrier of an audio broadcast, comprising:
 - identifying digital data associated with, but different than a set of broadcast data;
 - determining synchronization data corresponding to a broadcast event within the set of broadcast data;
 - assembling a data packet including the digital data and the synchronization data; and
 - transmitting the data packet to a receiver before the broadcast event, such that the broadcast data is presented with the digital data based on the synchronization data to a consumer of the broadcast data.
2. The method of claim 1, wherein the digital data corresponds to multimedia information for display on the receiver.
3. The method of claim 2 wherein the multimedia information may be stored on the receiver for recall by the consumer at a desired time.
4. The method of claim 1, wherein said transmitting further comprises transmitting the data packet in a digital format on

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at least one RF carrier of a central analog frequency on which the set of broadcast data is transmitted.

5. The method of claim 1 further comprising transmitting the set of broadcast data on a central analog frequency.

6. The method of claim 1 further comprising determining authentication data for allowing the receiver to authenticate the data packet.

7. The method of claim 6 wherein the assembling further comprises including the authentication data in the data packet for allowing the receiver to authenticate the source of the data packet.

8. The method of claim 6 wherein said authentication data includes at least one of a time stamp and nonce information.

9. The method of claim 1 wherein the assembling comprises encrypting the data packet with a public key.

10. The method of claim 1 further comprising determining a file size of the data packet.

11. The method of claim 10 wherein the transmitting further comprises selecting at least one RF carrier of a central analog frequency based on the file size.

12. The method of claim 11 wherein the transmitting further comprises transmitting the data packet on the selected RF carrier.

13. The method of claim 1 wherein said synchronization data comprises an indication that the digital data be stored by the receiver for recall by the consumer at a desired time.

14. The method of claim 4 wherein the central analog frequency corresponds to a single frequency on one of an amplitude modulated band and a frequency modulated band.

15. The method of claim 1 wherein said data packet comprises data in an extensible mark-up language format.

16. The method of claim 1, wherein the broadcast data is analog data.

17. The method of claim 1 wherein the broadcast data is digital data.

18. A computer comprising:
 a user interface; and
 a control system operatively coupled to the user interface and adapted to:
 identify a set of broadcast data that will be sent to a consumer through an RF channel as part of a one-way audio broadcast;
 identify digital data associated with, but different than the set of broadcast data;
 determine synchronization data corresponding to a broadcast event within the set of broadcast data;
 assemble a data packet comprising:
 the digital data; and
 the synchronization data; and
 transmit the data packet to a receiver before the broadcast event using an in-band, on-channel format, such that the set of broadcast data is presented with the digital data based on the synchronization data to the consumer of the broadcast data.

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